

Faster Binary Embeddings for Preserving Euclidean Distances

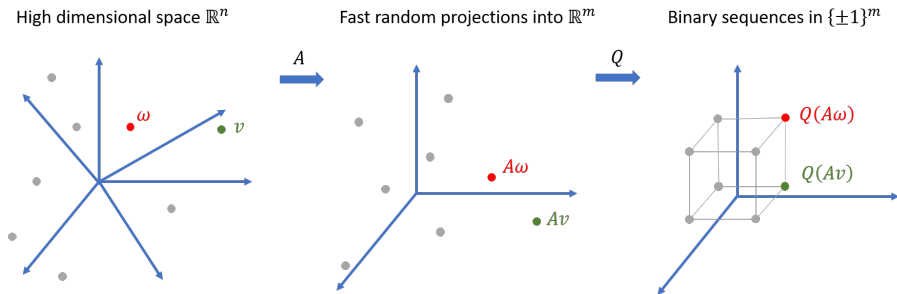
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Wednesday 24th March, 2021

Framework



Euclidean distance recovery: $d(Q(A\omega), Q(Av)) \approx \|\omega - v\|_2$

- Our method is

$$q_x := Q(Ax)$$

where Q is a stable $\Sigma\Delta$ quantization scheme, $A \in \mathbb{R}^{m \times n}$ is a sparse Gaussian matrix and $x \in \mathbb{R}^n$ is well-spread, i.e., those that are not sparse.

Sigma-Delta Quantization Schemes Q

For simplicity, we focus on the quantizer $Q : \mathbb{R}^m \rightarrow \{1, -1\}^m$. Specifically, $q = Q(x)$ such that for $i = 1, \dots, m$

$$\begin{cases} u_0 = 0, \\ q_i = \text{sign}(x_i + u_{i-1}), \\ u_i = u_{i-1} + x_i - q_i. \end{cases} \quad (1)$$

Algorithm 1: Fast Binary Embedding for Finite \mathcal{T}

Input: $\mathcal{T} = \{x^{(j)}\}_{j=1}^k \subseteq B_2^n(\kappa)$

▷ Data points in ℓ_2 ball

Generate $A \in \mathbb{R}^{m \times n}$

▷ Sparse Gaussian matrix A

for $j \leftarrow 1$ **to** k **do**

$z^{(j)} \leftarrow Ax^{(j)}$
 $q^{(j)} = Q(z^{(j)})$

▷ Stable $\Sigma\Delta$ quantizer Q as in (1)

Output: Binary sequences $\mathcal{B} = \{q^{(j)}\}_{j=1}^k \subseteq \{-1, 1\}^m$

ℓ_2 Distance Recovery

Algorithm 2: ℓ_2 Norm Distance Recovery

Input: $q^{(i)}, q^{(j)} \in \mathcal{B}$ \triangleright Binary sequences produced by Algorithm 1
 $y^{(i)} \leftarrow \tilde{V}q^{(i)}$ \triangleright Condense the components of q
 $y^{(j)} \leftarrow \tilde{V}q^{(j)}$

Output: $\|y^{(i)} - y^{(j)}\|_1$ \triangleright Approximation of $\|x^{(i)} - x^{(j)}\|_2$

Definition (Condensation operator)

Let p, r, λ be fixed positive integers such that $\lambda = r\tilde{\lambda} - r + 1$ for some integer $\tilde{\lambda}$. Let $m = \lambda p$ and v be a row vector in \mathbb{R}^λ whose entry v_j is the j -th coefficient of the polynomial $(1 + z + \dots + z^{\tilde{\lambda}-1})^r$. Define the normalized condensation operator $\tilde{V} \in \mathbb{R}^{p \times m}$ by

$$\tilde{V} = \frac{\sqrt{\pi/2}}{p\|v\|_2} \begin{bmatrix} v & & \\ & \dots & \\ & & v \end{bmatrix}.$$

Main Result

Theorem (Main result)

Let $\mathcal{T} \subseteq \mathbb{R}^n$ be a finite, appropriately scaled set with elements satisfying $\|x\|_\infty = O(n^{-1/2}\|x\|_2)$. If $m \gtrsim p := \Omega(\epsilon^{-2} \log(|\mathcal{T}|^2/\delta))$ and $r \geq 1$ is the integer order of Q , then with probability $1 - 2\delta$ on the draw of the sparse Gaussian matrix A , the following holds uniformly over all x, y in \mathcal{T} : Embedding x, y into $\{-1, 1\}^m$ using Algorithm 1, and estimating the associated distance between them using Algorithm 2 yields the error bound

$$\left| \|\tilde{V}(q_x - q_y)\|_1 - \|x - y\|_2 \right| \leq c \left(\frac{m}{p} \right)^{-r+1/2} + \epsilon \|x - y\|_2.$$

- The assumption that $\|x\|_\infty = O(n^{-1/2}\|x\|_2)$ is reasonable.
- The latter part in error bound is essentially proportional to $p^{-1/2}$.

Comparisons with Baselines

Method	Time	Space	Storage	Query Time
Toeplitz [1]	$O(n \log n)$	$O(n)$	$O(m)$	$O(m)$
Bilinear [2]	$O(n\sqrt{m})$	$O(\sqrt{mn})$	$O(m)$	$O(m)$
Circulant [3]	$O(n \log n)$	$O(n)$	$O(m)$	$O(m)$
BOE or PCE [4]	$O(n \log n)$	$O(n)$	$O(p \log_2 \lambda)$	$O(p \log_2^2 \lambda)$
Our Algorithm [5]	$O(m)$	$O(m)$	$O(p \log_2 \lambda)$	$O(p \log_2 \lambda)$

- “Time” is the time needed to embed a data point;
- “Space” is the space needed to store the embedding matrix;
- “Storage” contains the memory usage to store each encoded sequence;
- “Query time” is the time complexity of pairwise distance estimation.

Experiments on different Image Datasets

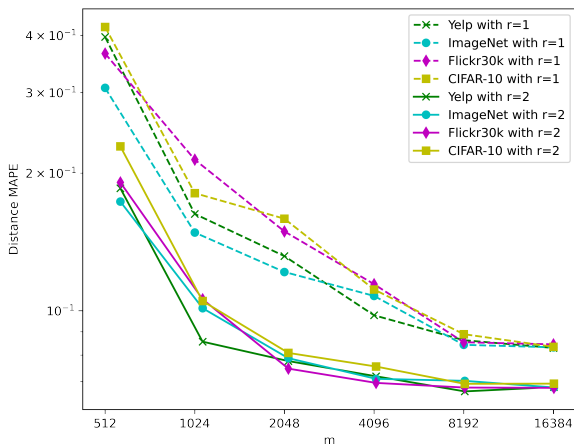


Figure: Plot of ℓ_2 distance reconstruction error on four datasets

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Thank You