# Scaling the Convex Barrier with Active Sets

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University of Oxford

Neural networks lack robustness (adversarial examples).

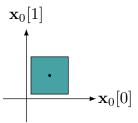
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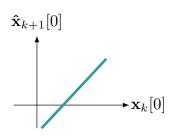
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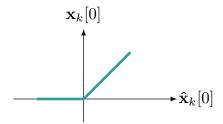
$$\mathbf{\hat{x}}_{k+1} = W_{k+1}\mathbf{x}_k + \mathbf{b}_{k+1}$$



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Verification as optimisation:

$$\mathbf{x}_k = \sigma\left(\mathbf{\hat{x}}_k\right)$$



[Bunel et al., 2018]

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#### Check the sign of:

$$\min_{\mathbf{x}, \hat{\mathbf{x}}} \hat{x}_n$$
s.t.  $\mathbf{x}_0 \in \mathcal{C}$ ,
$$\hat{\mathbf{x}}_{k+1} = W_{k+1}\mathbf{x}_k + \mathbf{b}_{k+1} \quad k \in \llbracket 0, n-1 \rrbracket,$$

$$\mathbf{x}_k = \sigma(\hat{\mathbf{x}}_k) \qquad \qquad k \in \llbracket 1, n-1 \rrbracket.$$

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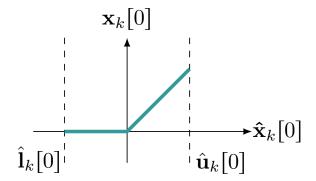
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NP-HARD

### Approximate $\min \hat{x}_n$ via a lower bound:

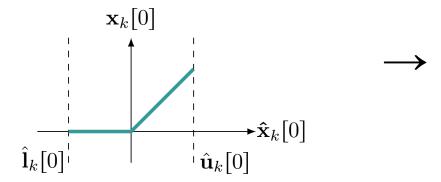
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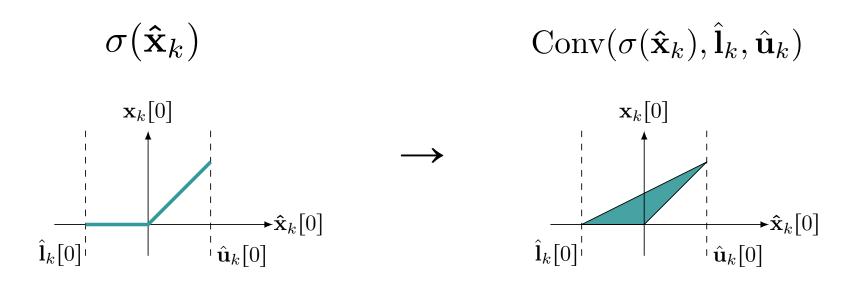


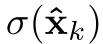
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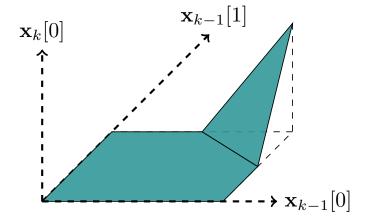
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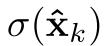


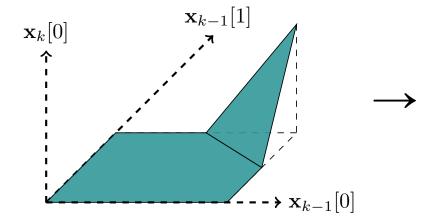
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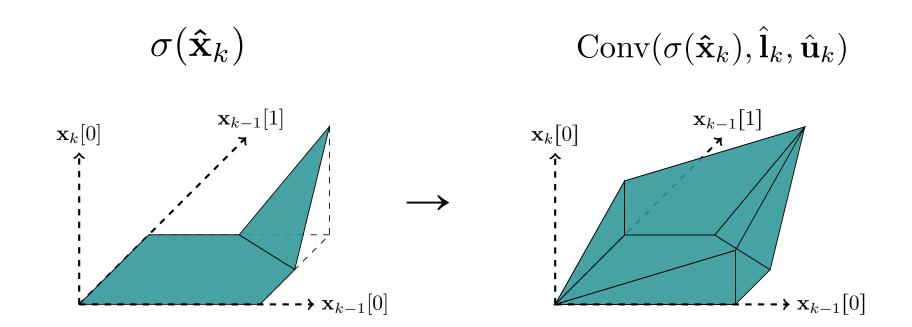


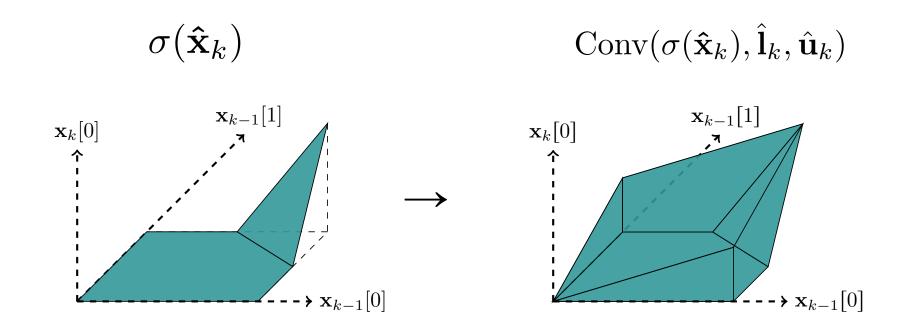






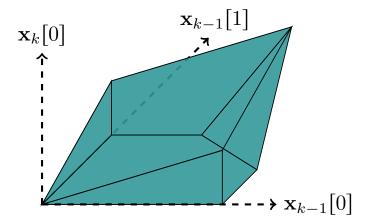




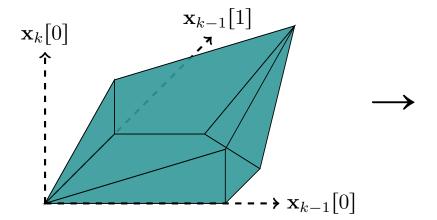


Loose.

$$\operatorname{Conv}(\sigma(\mathbf{\hat{x}}_k), \hat{\mathbf{l}}_k, \hat{\mathbf{u}}_k)$$



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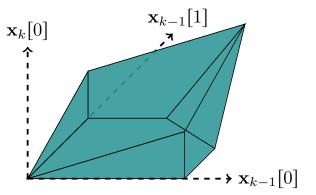
$$\overset{\mathbf{x}_{k}[0]}{\longrightarrow} \overset{\mathbf{x}_{k-1}[1]}{\longrightarrow} \overset{\mathbf{x}_{k}[0]}{\longrightarrow} \overset{\mathbf{x}_{k-1}[0]}{\longrightarrow} \overset{\mathbf{x}_{k$$

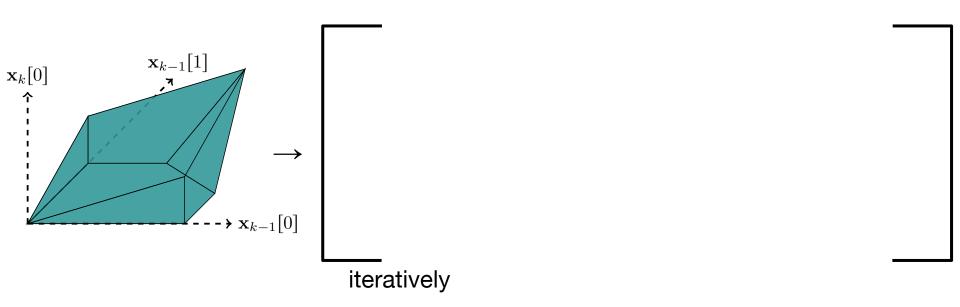
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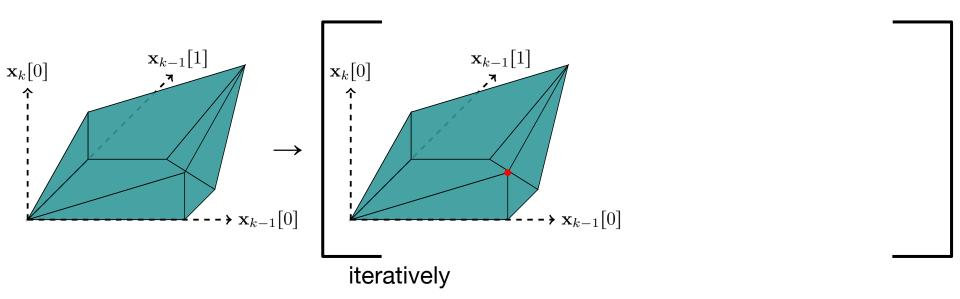
$$\mathbf{x}_{k}[0] \qquad \mathbf{x}_{k-1}[1] \qquad \mathbf{x}_{k}[0] \qquad \mathbf{x}_{k-1}[1] \qquad \mathbf{x}_{k-1}[0]$$

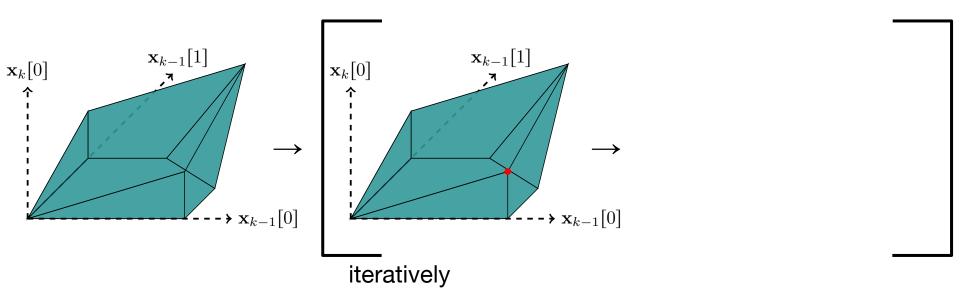
#### Exponentially many constraints.

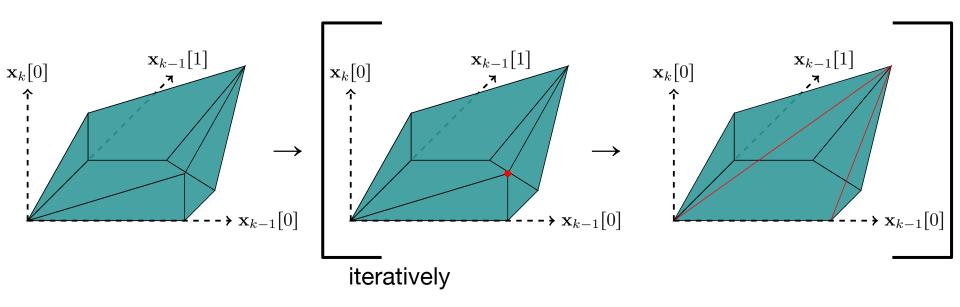
Figures modified from [Anderson et al., 2020]



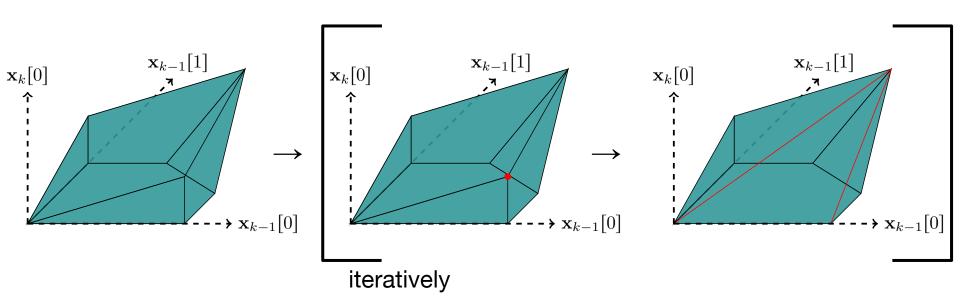








Most violated constraint from  $Conv(\sigma(W_k\mathbf{x}_{k-1} + \mathbf{b}_k), \hat{\mathbf{l}}_k, \hat{\mathbf{u}}_k)$  at any point found in *linear-time*.



Off-the-shelf LP solvers scale poorly with NN size.

[Bunel et al., 2020]

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An efficient solver for the tight relaxation needs:

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1. anytime property

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$$\max_{\substack{\mathbf{x}_{k-1}[1] \\ \mathbf{x}_{k}[0] \\ \vdots}} \left\{ \min_{\substack{(\boldsymbol{\alpha}, \boldsymbol{\beta}_{\mathcal{A}}) \geq 0}} \left\{ \min_{\substack{(\mathbf{x}, \mathbf{z}) \in \mathcal{D}}} \mathcal{L}(\mathbf{x}, \mathbf{z}, \boldsymbol{\alpha}, \boldsymbol{\beta}_{\mathcal{A}}) \right\}$$

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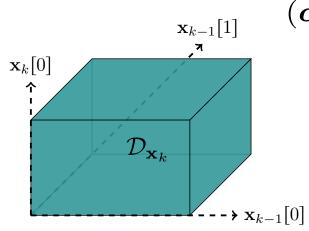
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Exponentially many variables.

$$\max_{(\boldsymbol{\alpha},\boldsymbol{\beta}_{\mathcal{A}})\geq 0} \left\{ \min_{(\mathbf{x},\mathbf{z})\in\mathcal{D}} \mathcal{L}(\mathbf{x},\mathbf{z},\boldsymbol{\alpha},\boldsymbol{\beta}_{\mathcal{A}}) \right\}$$



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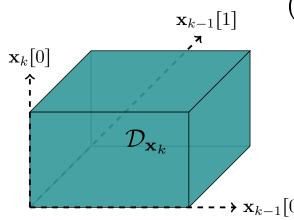
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#### 2. sparsity

Selection criterion for active set B?

$$\max_{(\boldsymbol{\alpha},\boldsymbol{\beta}_{\mathcal{B}}) \geqslant 0} \left\{ \min_{(\mathbf{x},\mathbf{z}) \in \mathcal{D}} \mathcal{L}_{\mathcal{B}}(\mathbf{x},\mathbf{z},\boldsymbol{\alpha},\boldsymbol{\beta}_{\mathcal{B}}) \right\}$$



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3. tightness

Start from Planet dual;

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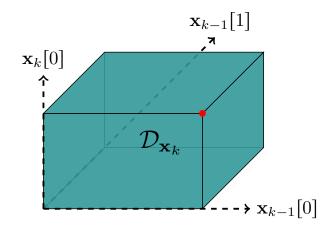
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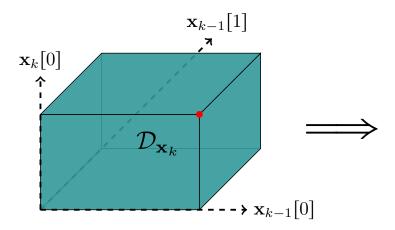
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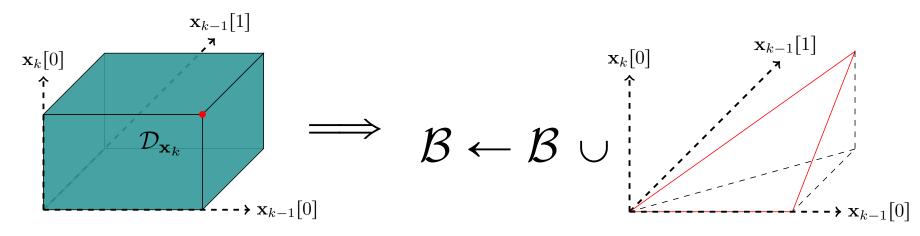
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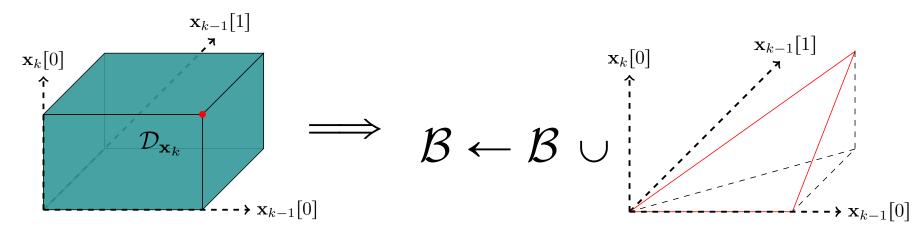
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### **Active Set**

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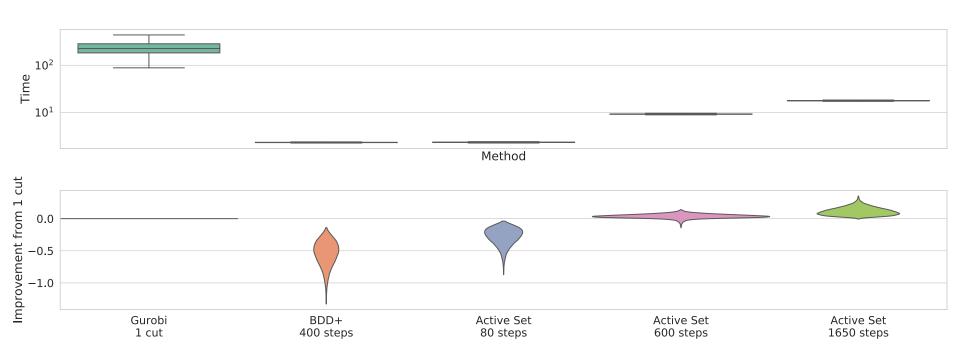
# **Experiments: Baselines**

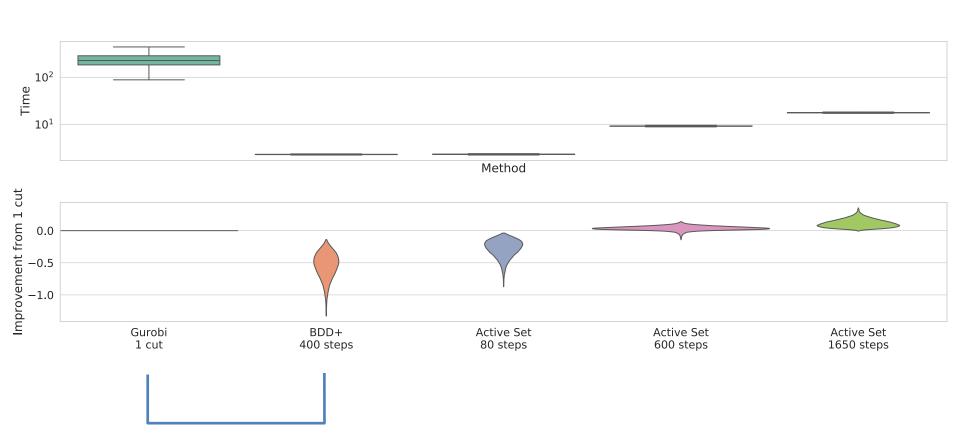
• **BDD**+: Efficient dual solver for Planet relaxation from previous work.

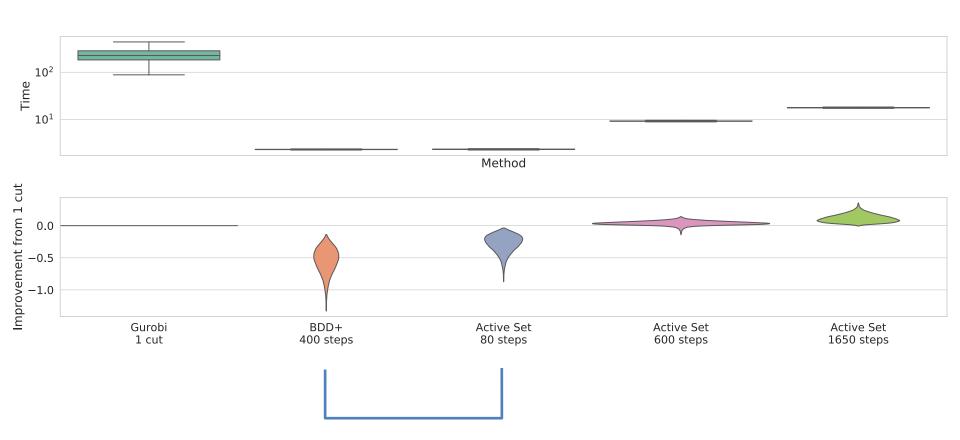
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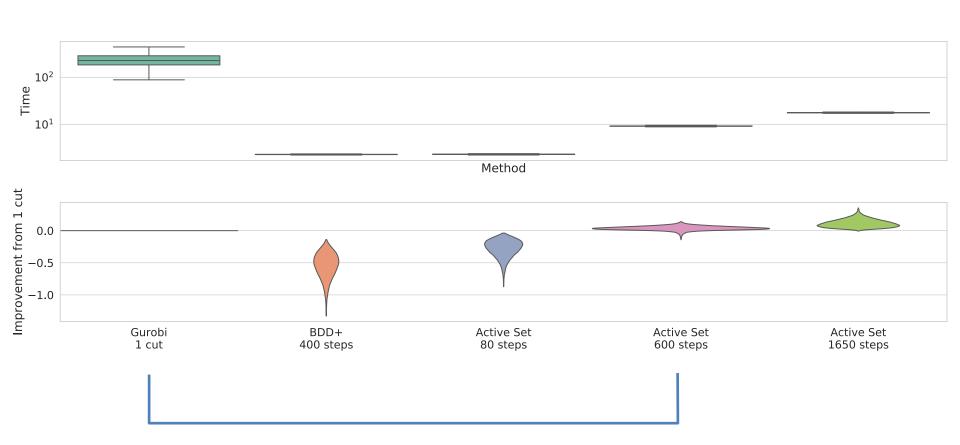
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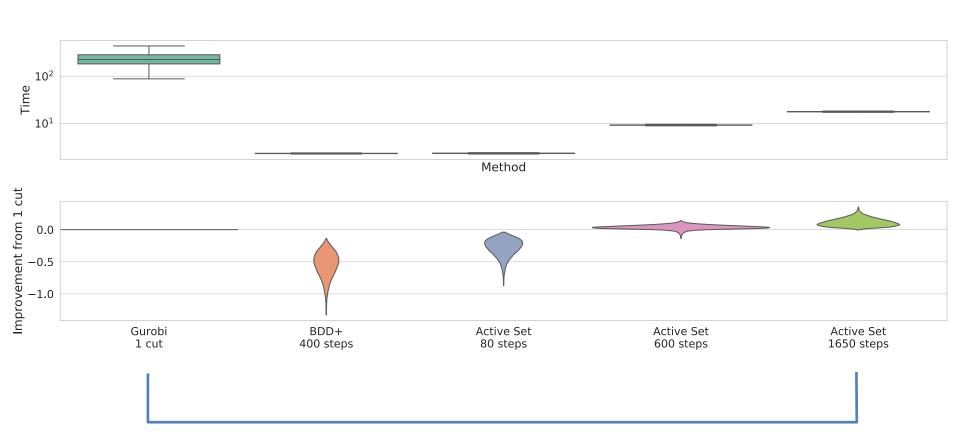
Gurobi: Primal cutting plane algorithm with one cut.











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Customised dual solver for tight ReLU relaxation: <u>a new convex barrier</u>?

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Follow-up work: more memory efficient solver.

[arXiv: Scaling the Convex Barrier with Sparse Dual Algorithms]