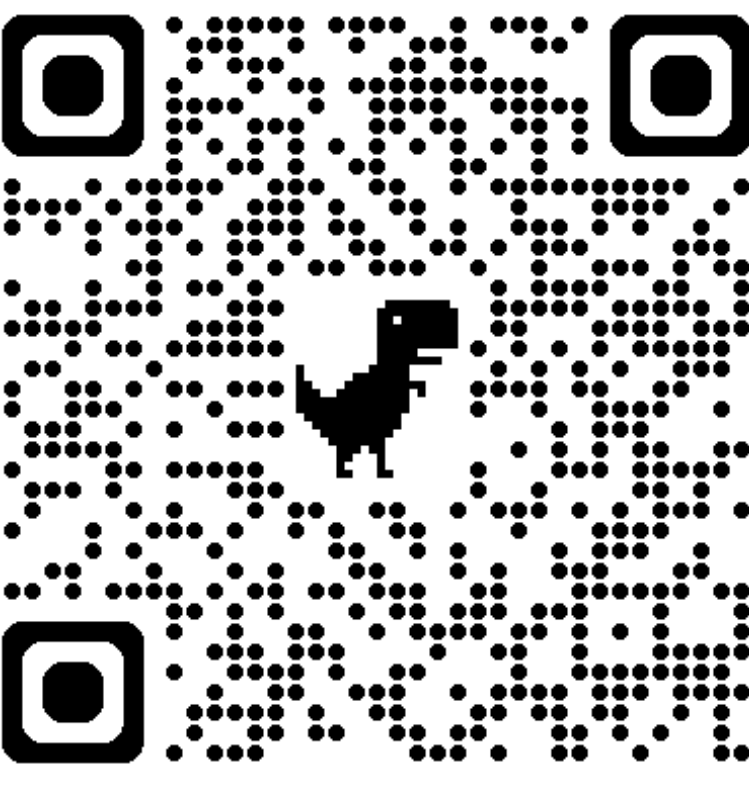


Permute-and-Flip: An optimally stable and watermarkable decoder for LLMs



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Existing watermarks work with the standard decoder:
softmax(logits)

Softmax sampling: $y_t \sim p(y) = \frac{e^{u(y|x, y_{1:t-1})/T}}{\sum_{\tilde{y}} e^{u(\tilde{y}|x, y_{1:t-1})/T}}$

Is softmax(logits) the optimal choice?



Can we benefit from co-designing the decoder and watermarking scheme?

PF Decoding

Permute-and-Flip Decoding

Sampling

$$y_t = \arg \max_{y \in \mathcal{V}} \frac{u_t(y)}{T} + G_t(y)$$

$G_t(y) \sim \text{Gumbel}(0, 1) \text{ i.i.d.}$

PF Decoding

$$y_t = \arg \max_{y \in \mathcal{V}} \frac{u_t(y)}{T} + E_t(y)$$

$E_t(y) \sim \text{Exponential}(1) \text{ i.i.d.}$

PF Watermark

Gumbel Watermark [Aaronson, 2023]

$$y_t = \arg \max_{y \in \mathcal{V}} \frac{u_t(y)}{T} + G_t(y)$$

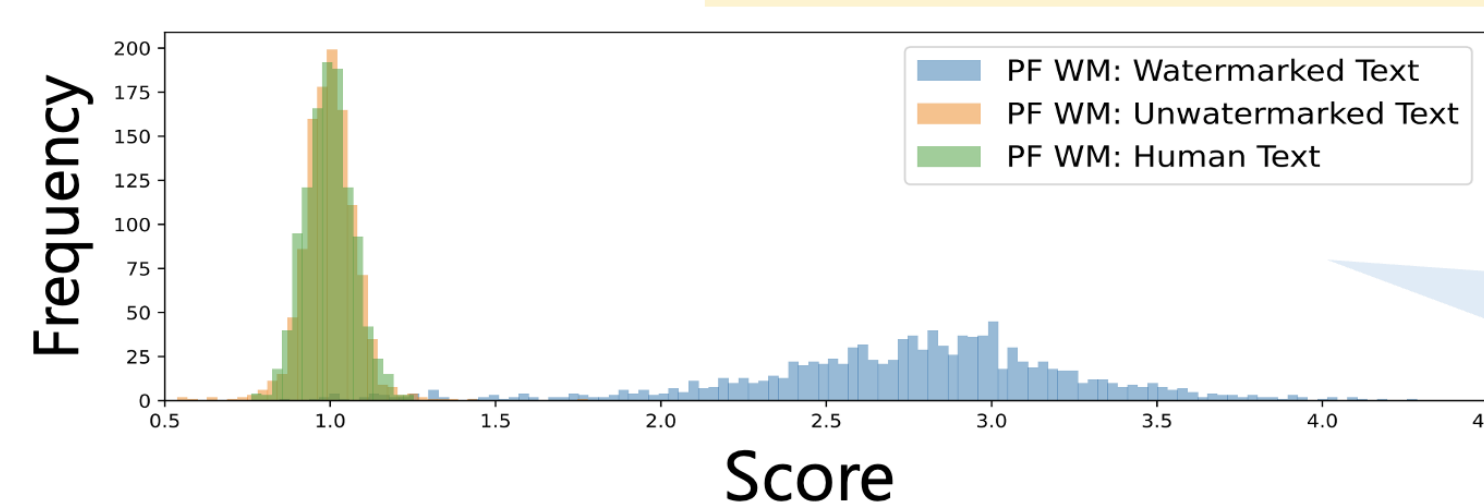
$G_t(y) \sim \text{Gumbel}(0, 1) \text{ i.i.d.}$

PF Watermark [Ours]

$$y_t = \arg \max_{y \in \mathcal{V}} \frac{u_t(y)}{T} + E_t(y)$$

$E_t(y) \sim \text{Exponential}(1) \text{ i.i.d.}$

Make them Pseudo-random!



A Clear Gap
↓
Watermark Detectability

Stability

Definition 2.1 (Stability). We say a decoding algorithm $\mathcal{A}$ is L -stable if for any prompt x , prefix $y_{\leq t}$, and for any perturbed $\tilde{u}$ such that $\|\tilde{u} - u\|_{\infty} \leq \delta$, the log-probability ratio satisfies

$$\left| \log \left\{ \frac{p_{\mathcal{A}(\tilde{u}(\cdot|x, y_{\leq t}))}(y)}{p_{\mathcal{A}(u(\cdot|x, y_{\leq t}))}(y)} \right\} \right| \leq L\delta \quad \forall y \in \mathcal{V}. \quad \begin{array}{l} \blacksquare \text{ Avoid catastrophic failure} \\ \blacksquare \text{ Implies diversity} \end{array}$$

Guarantees

Theorem 3.1. Let the logits function be u and $u^* = \max_{y \in \mathcal{V}} u(y)$. Let $\text{PF}(u)$ be the distribution of PF-sampling, and $\text{Softmax}(u)$ be the distribution in equation 1, both with temperature parameter T . The following statements are true.

1. (**Same stability**) PF-Sampling is $(2/T)$ -stable.
2. (**Nearly greedy**) PF-sampling obeys,

$$\mathbb{E}_{y \sim \text{PF}(u)}[u(y)] \geq u^* - T \log |\mathcal{V}|.$$

3. (**“Never worse”**) For the same T , PF-sampling is never worse than Softmax-sampling.

$$\mathbb{E}_{y \sim \text{PF}(u)}[u(y)] \geq \mathbb{E}_{y \sim \text{Softmax}(u)}[u(y)]$$

4. (**“Up to 2x better”**) There exists logits u such that PF-sampling is 2x smaller in terms of suboptimality,

$$\mathbb{E}_{y \sim \text{PF}(u)}[u^* - u(y)] \leq \frac{1}{2} \mathbb{E}_{y \sim \text{Softmax}(u)}[u^* - u(y)].$$

5. (**Optimal stability-perplexity tradeoff**) For any decoder P that is $2/T$ -stable, if there exists u such that

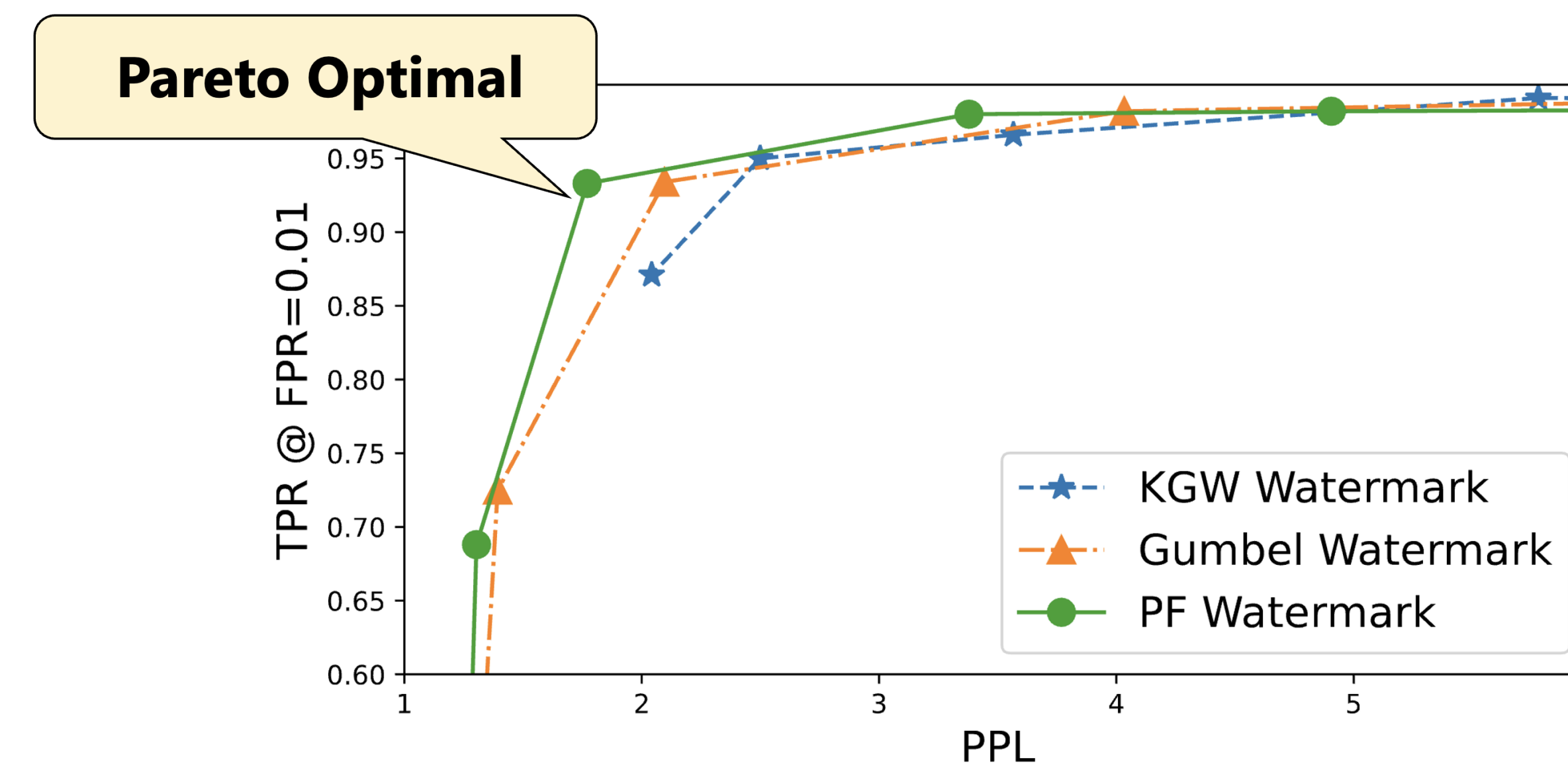
$$\mathbb{E}_{y \sim P(u)}[u(y)] > \mathbb{E}_{y \sim \text{PF}(u)}[u(y)]$$

then there must be another $\tilde{u}$ such that

$$\mathbb{E}_{y \sim P(\tilde{u})}[\tilde{u}(y)] < \mathbb{E}_{y \sim \text{PF}(\tilde{u})}[\tilde{u}(y)].$$

Results

Better Detectability-Perplexity Tradeoff



Method	AUC↑	TPR↑	PPL1↓	PPL2↓	Seq-rep-5↓	MAUVE↑
C4, T=1.0, Llama2-7B						
Greedy	-	-	1.14 _{0.01}	1.24 _{0.03}	0.56	0.05
Sampling	-	-	12.47 _{0.32}	15.31 _{0.41}	0.02	0.98
PF	-	-	8.94 _{0.20}	10.75 _{0.25}	0.03	0.90
KGW WM	0.989	0.991	16.62 _{0.38}	20.62 _{0.49}	0.01	1.00
Gumbel WM	0.997	0.988	11.41 _{0.27}	14.12 _{0.36}	0.04	0.93
PF WM	0.995	0.984	8.33 _{0.20}	10.28 _{0.29}	0.05	0.99