

# Tight Lower Bounds under Asymmetric High-Order Hölder Smoothness and Uniform Convexity



**Cedar Site Bai**  
bai123@purdue.edu



**Brian Bullins**  
bbullins@purdue.edu

## I Oracle Complexity / Convergence Rate

### Oracle Complexity

- Objective:  $\min_x f(x)$
- # times an algorithm accesses an oracle to reach an  $\epsilon$ -approximate solution

$$\text{e.g., } f(x_T) - f(x^*) \leq \epsilon$$

## I Oracle Complexity / Convergence Rate

### Oracle Complexity

- Objective:  $\min_x f(x)$
- # times an algorithm accesses an oracle to reach an  $\epsilon$ -approximate solution

$$\text{e.g., } f(x_T) - f(x^*) \leq \epsilon$$

- First-order Oracle: gradient  $\nabla f(x)$
- First-order oracle complexity: # times an algorithm computes gradient

## I Oracle Complexity / Convergence Rate

### Oracle Complexity

- Objective:  $\min_x f(x)$
- # times an algorithm accesses an oracle to reach an  $\epsilon$ -approximate solution  
 e.g.,  $f(x_T) - f(x^*) \leq \epsilon$
- First-order Oracle: gradient  $\nabla f(x)$
- First-order oracle complexity: # times an algorithm computes gradient
- Oracle complexity of gradient descent for convex and smooth functions:  $T \in \Theta(\epsilon^{-1})$

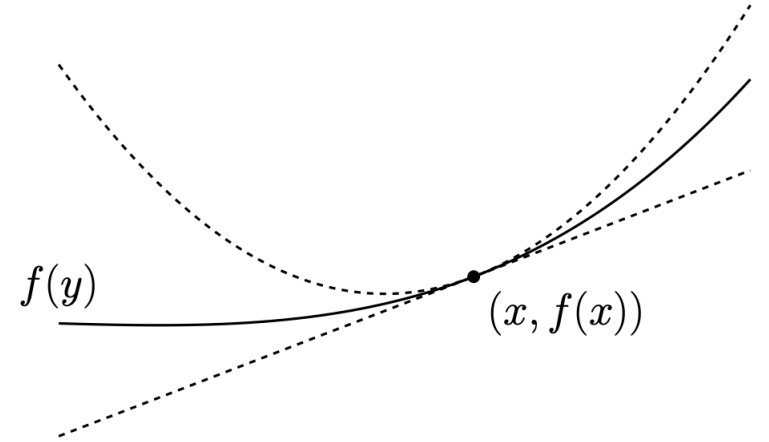
### Convergence Rate

- How fast the approximation error decays with # iterations
- Convergence rate of gradient descent (convex & smooth):  $f(x_T) - f(x^*) \leq \epsilon \in \Theta\left(\frac{1}{T}\right)$

## | Smoothness

$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, L > 0, \|\nabla f(\mathbf{x}) - \nabla f(\mathbf{y})\| \leq L \|\mathbf{x} - \mathbf{y}\|$$

$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, L > 0, f(\mathbf{y}) - f(\mathbf{x}) - \langle \nabla f(\mathbf{x}), \mathbf{y} - \mathbf{x} \rangle \leq \frac{L}{2} \|\mathbf{x} - \mathbf{y}\|^2$$

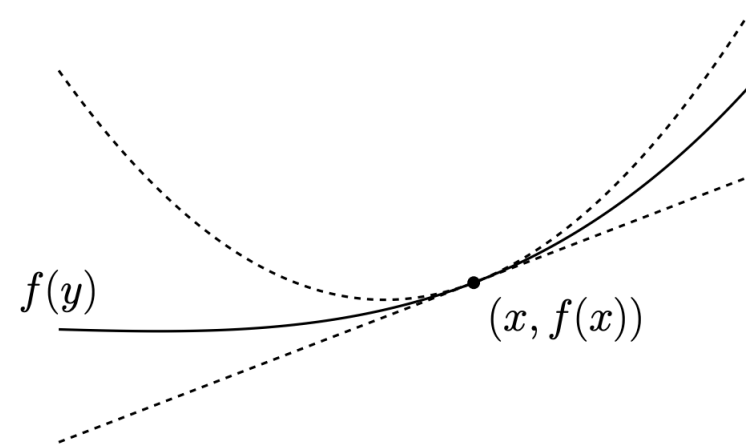


## | Smoothness

$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, L > 0, \|\nabla f(\mathbf{x}) - \nabla f(\mathbf{y})\| \leq L \|\mathbf{x} - \mathbf{y}\|$$

## | Hölder Smoothness

$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, H > 0, \nu \in (0, 1], \|\nabla f(\mathbf{x}) - \nabla f(\mathbf{y})\| \leq H \|\mathbf{x} - \mathbf{y}\|^\nu$$



## | Smoothness

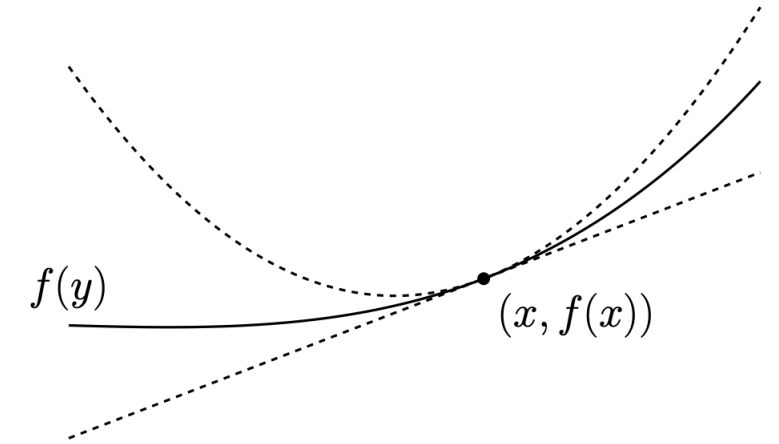
$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, L > 0, \|\nabla f(\mathbf{x}) - \nabla f(\mathbf{y})\| \leq L \|\mathbf{x} - \mathbf{y}\|$$

## | Hölder Smoothness

$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, H > 0, \nu \in (0, 1], \|\nabla f(\mathbf{x}) - \nabla f(\mathbf{y})\| \leq H \|\mathbf{x} - \mathbf{y}\|^\nu$$

## | High-order Smoothness

$$p \in \mathbb{Z}^+, \forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, L > 0, \|\nabla^p f(\mathbf{x}) - \nabla^p f(\mathbf{y})\| \leq L \|\mathbf{x} - \mathbf{y}\|$$



## | Smoothness

$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, L > 0, \|\nabla f(\mathbf{x}) - \nabla f(\mathbf{y})\| \leq L \|\mathbf{x} - \mathbf{y}\|$$

## | Hölder Smoothness

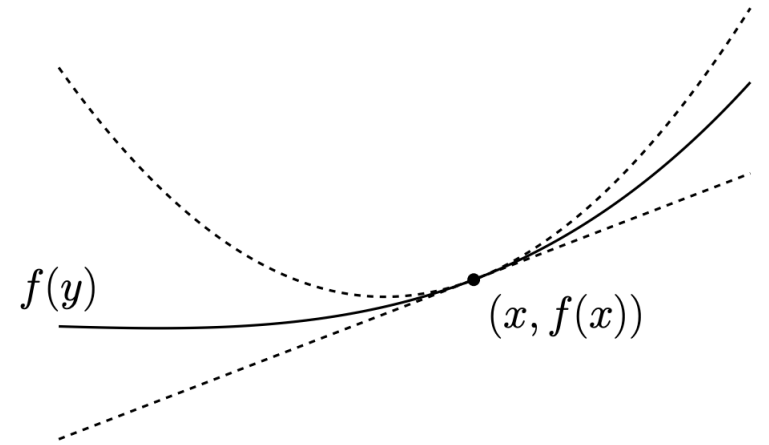
$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, H > 0, \nu \in (0, 1], \|\nabla f(\mathbf{x}) - \nabla f(\mathbf{y})\| \leq H \|\mathbf{x} - \mathbf{y}\|^\nu$$

## | High-order Smoothness

$$p \in \mathbb{Z}^+, \forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, L > 0, \|\nabla^p f(\mathbf{x}) - \nabla^p f(\mathbf{y})\| \leq L \|\mathbf{x} - \mathbf{y}\|$$

## | High-order Hölder Smoothness

$$p \in \mathbb{Z}^+, \forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, H > 0, \nu \in (0, 1], \|\nabla^p f(\mathbf{x}) - \nabla^p f(\mathbf{y})\| \leq H \|\mathbf{x} - \mathbf{y}\|^\nu$$



# I Convexity

$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, \langle \nabla f(\mathbf{y}) - \nabla f(\mathbf{x}), \mathbf{y} - \mathbf{x} \rangle \geq 0$$

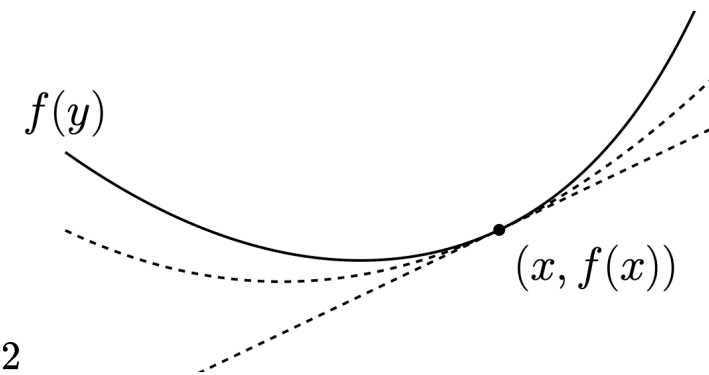
$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, f(\mathbf{y}) - f(\mathbf{x}) - \langle \nabla f(\mathbf{x}), \mathbf{y} - \mathbf{x} \rangle \geq 0$$

## | Convexity

$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, \langle \nabla f(\mathbf{y}) - \nabla f(\mathbf{x}), \mathbf{y} - \mathbf{x} \rangle \geq 0$$

## | Strong Convexity

$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, \sigma > 0, \langle \nabla f(\mathbf{y}) - \nabla f(\mathbf{x}), \mathbf{y} - \mathbf{x} \rangle \geq \sigma \|\mathbf{x} - \mathbf{y}\|^2$$



## I Convexity

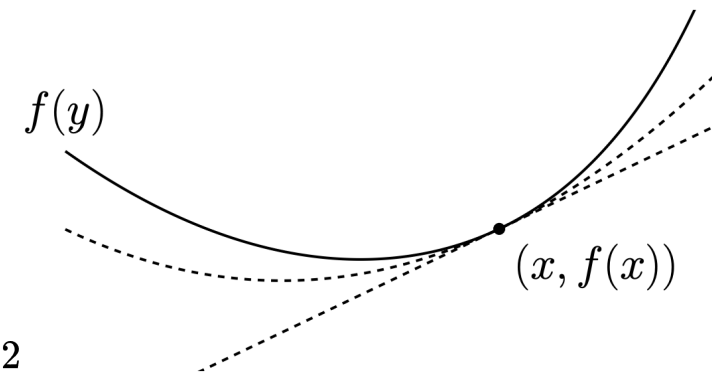
$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, \langle \nabla f(\mathbf{y}) - \nabla f(\mathbf{x}), \mathbf{y} - \mathbf{x} \rangle \geq 0$$

## I Strong Convexity

$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, \sigma > 0, \langle \nabla f(\mathbf{y}) - \nabla f(\mathbf{x}), \mathbf{y} - \mathbf{x} \rangle \geq \sigma \|\mathbf{x} - \mathbf{y}\|^2$$

## I Uniform Convexity

$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, \sigma > 0, q \geq 2, \langle \nabla f(\mathbf{y}) - \nabla f(\mathbf{x}), \mathbf{y} - \mathbf{x} \rangle \geq \sigma \|\mathbf{x} - \mathbf{y}\|^q$$



## I Motivating Example

### Subproblem of Cubic-regularized Newton Method

Let  $M$  be a positive parameter. Define a modified Newton step using the following *cubic regularization* of quadratic approximation of function  $f(x)$ :

$$T_M(x) \in \operatorname{Arg\,min}_y \left[ \langle f'(x), y-x \rangle + \frac{1}{2} \langle f''(x)(y-x), y-x \rangle + \frac{M}{6} \|y-x\|^3 \right], \quad (2.4)$$

**Cubic regularization of Newton method**

**Initialization:** Choose  $x_0 \in R^d$ .

**Iteration  $k$ , ( $k \geq 0$ ):**

1. Find  $M_k \in [L_0, 2L]$  such that  $f(T_{M_k}(x_k)) \leq \bar{f}_{M_k}(x_k)$ .
2. Set  $x_{k+1} = T_{M_k}(x_k)$ .

## I Motivating Example

### Subproblem of Cubic-regularized Newton Method

$$F(y) := \langle \nabla f(x_k), y - x_k \rangle + \frac{1}{2} \langle \nabla^2 f(x_k)(y - x_k), y - x_k \rangle + \frac{M}{6} \|y - x_k\|^3$$

$$x_{k+1} \in \arg \min_y F(y)$$

**Cubic regularization of Newton method**

**Initialization:** Choose  $x_0 \in \mathbb{R}^d$ .

**Iteration  $k$ , ( $k \geq 0$ ):**

1. Find  $M_k \in [L_0, 2L]$  such that  
 $f(T_{M_k}(x_k)) \leq \bar{f}_{M_k}(x_k)$ .
2. Set  $x_{k+1} = T_{M_k}(x_k)$ .

## I Motivating Example

### Subproblem of Cubic-regularized Newton Method

$$F(y) := \langle \nabla f(x_k), y - x_k \rangle + \frac{1}{2} \langle \nabla^2 f(x_k)(y - x_k), y - x_k \rangle + \frac{M}{6} \|y - x_k\|^3$$

- Uniformly convex with degree  $q=3$  and parameter  $\sigma$ :

$$F(y) - F(x) - \langle \nabla F(x), y - x \rangle \geq \frac{\sigma}{q} \|y - x\|^q$$

- Second-order smooth / Lipschitz Hessian:

$$\|\nabla^2 F(y) - \nabla^2 F(x)\| \leq M \|y - x\|$$

## I Motivating Example

### Subproblem of Cubic-regularized Newton Method

$$F(y) := \langle \nabla f(x_k), y - x_k \rangle + \frac{1}{2} \langle \nabla^2 f(x_k)(y - x_k), y - x_k \rangle + \frac{M}{6} \|y - x_k\|^3$$

- Uniformly convex with degree  $q=3$  and parameter  $\sigma$ :

$$F(y) - F(x) - \langle \nabla F(x), y - x \rangle \geq \frac{\sigma}{q} \|y - x\|^q$$

- Second-order smooth / Lipschitz Hessian:

$$\|\nabla^2 F(y) - \nabla^2 F(x)\| \leq M \|y - x\|$$

- Ideally, solve with optimal second-order methods [Gasnikov et. al., 2019]

$$p = 2, q = 3, \nu = 1, q = p + \nu, \implies \mathcal{O} \left( \left( \frac{M}{\sigma} \right)^{\frac{2}{7}} \log \left( \frac{1}{\epsilon} \right) \right)$$

## I Motivating Example

### Subproblem of Cubic-regularized Newton Method

$$F(y) := \langle \nabla f(x_k), y - x_k \rangle + \frac{1}{2} \langle \nabla^2 f(x_k)(y - x_k), y - x_k \rangle + \frac{M}{6} \|y - x_k\|^3$$

- Ideally, solve with optimal second-order methods [Gasnikov et. al., 2019]

$$p = 2, q = 3, \nu = 1, q = p + \nu \implies \mathcal{O} \left( \left( \frac{M}{\sigma} \right)^{\frac{2}{7}} \log \left( \frac{1}{\epsilon} \right) \right)$$

- Computing Hessian is expensive:  $\mathcal{O}(d^2)$
- Use first-order methods, e.g., (accelerated) gradient descent, instead:  $\mathcal{O}(d)$

## I Motivating Example

### Subproblem of Cubic-regularized Newton Method

$$F(y) := \langle \nabla f(x_k), y - x_k \rangle + \frac{1}{2} \langle \nabla^2 f(x_k)(y - x_k), y - x_k \rangle + \frac{M}{6} \|y - x_k\|^3$$

- Ideally, solve with optimal second-order methods [Gasnikov et. al., 2019]

$$p = 2, q = 3, \nu = 1, \boxed{q = p + \nu} \implies \mathcal{O} \left( \left( \frac{M}{\sigma} \right)^{\frac{2}{7}} \log \left( \frac{1}{\epsilon} \right) \right)$$

- Computing Hessian is expensive:  $\mathcal{O}(d^2)$
- Use first-order methods, e.g., (accelerated) gradient descent, instead:  $\mathcal{O}(d)$
- **Order of smoothness** limited by **order of oracle**, e.g., H-smooth (within some domain)

$$p = 1, q = 3, \nu = 1, \boxed{q > p + \nu} \implies \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{1}{2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{1}{6}} \right) \quad [\text{Roulet \& d'Aspremont, 2017; Song et. al. 2021}]$$

# | Uniform Convexity and High-Order Hölder Smoothness

- Uniformly Convexity

$$\forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, \sigma > 0, q \geq 2, \langle \nabla f(\mathbf{y}) - \nabla f(\mathbf{x}), \mathbf{y} - \mathbf{x} \rangle \geq \sigma \|\mathbf{x} - \mathbf{y}\|^q$$

- High-order Hölder-smoothness

$$p \in \mathbb{Z}^+, \forall \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, H > 0, \nu \in (0, 1], \|\nabla^p f(\mathbf{x}) - \nabla^p f(\mathbf{y})\| \leq H \|\mathbf{x} - \mathbf{y}\|^\nu$$

- Question:

What's the  **$p^{\text{th}}$ -order oracle complexity** for **different combinations of  $p, q, \nu$** ?

# I Upper Bounds

Upper Bounds [Song et. al. 2021]

---


$$q > p + \nu \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+\nu)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-\nu)}{q(3(p+\nu)-2)}} \right)$$


---

$$q = p + \nu \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+\nu)-2}} \log \left( \frac{1}{\epsilon} \right) \right)$$


---

$$q < p + \nu \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+\nu)-2}} + \log \log \left( \left( \frac{\sigma^{p+\nu}}{H^q} \right)^{\frac{1}{p+\nu-q}} \frac{1}{\epsilon} \right) \right)$$


---

# I Upper Bounds

## Upper Bounds [Song et. al. 2021]

$$q > p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-v)}{q(3(p+v)-2)}} \right)$$

$p = 1$  [Roulet & d'Aspremont, 2017]

$$q = p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \log \left( \frac{1}{\epsilon} \right) \right)$$

[Gasnikov et. al., 2019]

$$q < p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} + \log \log \left( \left( \frac{\sigma^{p+v}}{H^q} \right)^{\frac{1}{p+v-q}} \frac{1}{\epsilon} \right) \right)$$

$q = 2, p = 2, v = 1$  [Arjevani et al., 2019]

# Lower Bounds

Upper Bounds [Song et. al. 2021]

**Lower Bounds [Our Work]**

$$q > p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-v)}{q(3(p+v)-2)}} \right)$$

$p = 1$  [Roulet & d'Aspremont, 2017]

$$\Omega \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-v)}{q(3(p+v)-2)}} \right)$$

$$q = p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \log \left( \frac{1}{\epsilon} \right) \right)$$

[Gasnikov et. al., 2019]

Future Work

$$q < p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} + \log \log \left( \left( \frac{\sigma^{p+v}}{H^q} \right)^{\frac{1}{p+v-q}} \frac{1}{\epsilon} \right) \right) \quad \Omega \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} + \log \log \left( \left( \frac{\sigma^{p+v}}{H^q} \right)^{\frac{1}{p+v-q}} \frac{1}{\epsilon} \right) \right)$$

$q = 2, p = 2, v = 1$  [Arjevani et al., 2019]

# Lower Bounds

## Upper Bounds [Song et. al. 2021]

$$q > p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-v)}{q(3(p+v)-2)}} \right)$$

$p = 1$  [Roulet & d'Aspremont, 2017]

$$q = p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \log \left( \frac{1}{\epsilon} \right) \right)$$

[Gasnikov et. al., 2019]

$$q < p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} + \log \log \left( \left( \frac{\sigma^{p+v}}{H^q} \right)^{\frac{1}{p+v-q}} \frac{1}{\epsilon} \right) \right)$$

$q = 2, p = 2, v = 1$  [Arjevani et al., 2019]

## Lower Bounds [Our Work]

$$\Omega \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-v)}{q(3(p+v)-2)}} \right)$$

$p = 1$  [Thomsen & Doikov, 2024]

Future Work

$$\Omega \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} + \log \log \left( \left( \frac{\sigma^{p+v}}{H^q} \right)^{\frac{1}{p+v-q}} \frac{1}{\epsilon} \right) \right)$$

$q = 2, p = 2, v = 1$  [Arjevani et al., 2019]  
 $v = 1$  [Kornowski & Shamir 2020]

## | **Hard Function** Construction for Case 1: $q > p + \nu$

- Start with a sequence of non-smooth functions

$$g_t(\mathbf{x}) = \max_{1 \leq k \leq t} r_k(\mathbf{x}) \quad \text{where} \quad r_k(\mathbf{x}) = \xi_k \langle \mathbf{e}_{\alpha(k)}, \mathbf{x} \rangle - (k-1)\delta, \forall k \in [T]$$

$\xi_k \in \{-1, 1\}$ ,  $\mathbf{e}$  : standard basis,  $\alpha$  : permutation of  $[T]$ ,  $\delta > 0$ .

# I Hard Function Construction for Case 1: $q > p + \nu$

- Start with a sequence of non-smooth functions

$$g_t(\mathbf{x}) = \max_{1 \leq k \leq t} r_k(\mathbf{x}) \quad \text{where} \quad r_k(\mathbf{x}) = \xi_k \langle \mathbf{e}_{\alpha(k)}, \mathbf{x} \rangle - (k-1)\delta, \forall k \in [T]$$

$\xi_k \in \{-1, 1\}$ ,  $\mathbf{e}$  : standard basis,  $\alpha$  : permutation of  $[T]$ ,  $\delta > 0$ .

- Smooth the functions with **a smoothing operator**, making it  $p^{\text{th}}$ -order **H-smooth**

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

# Hard Function Construction for Case 1: $q > p + \nu$

- Start with a sequence of non-smooth functions

$$g_t(\mathbf{x}) = \max_{1 \leq k \leq t} r_k(\mathbf{x}) \quad \text{where} \quad r_k(\mathbf{x}) = \xi_k \langle \mathbf{e}_{\alpha(k)}, \mathbf{x} \rangle - (k-1)\delta, \forall k \in [T]$$

$\xi_k \in \{-1, 1\}$ ,  $\mathbf{e}$  : standard basis,  $\alpha$  : permutation of  $[T]$ ,  $\delta > 0$ .

- **Smooth the functions** with a smoothing operator, making it  $p^{\text{th}}$ -order **H-smooth**

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

- Add regularization, making it  $q^{\text{th}}$ -order uniformly convex

$$F_t(\mathbf{x}) = \beta G_t(\mathbf{x}) + \frac{\sigma}{q} \|\mathbf{x}\|^q, \quad \mathbf{x} \in \mathcal{Q}, \beta > 0 \quad F(\mathbf{x}) = F_T(\mathbf{x}).$$

# I Hard Function Construction for Case 1: $q > p + \nu$

- Start with a sequence of non-smooth functions

$$g_t(\mathbf{x}) = \max_{1 \leq k \leq t} r_k(\mathbf{x}) \quad \text{where} \quad r_k(\mathbf{x}) = \xi_k \langle \mathbf{e}_{\alpha(k)}, \mathbf{x} \rangle - (k-1)\delta, \forall k \in [T]$$

$\xi_k \in \{-1, 1\}$ ,  $\mathbf{e}$  : standard basis,  $\alpha$  : permutation of  $[T]$ ,  $\delta > 0$ .

- Smooth the functions with a smoothing operator, making it  $p^{\text{th}}$ -order **H-smooth**

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

- Add regularization, making it  $q^{\text{th}}$ -order uniformly convex

$$F_t(\mathbf{x}) = \beta G_t(\mathbf{x}) + \frac{\sigma}{q} \|\mathbf{x}\|^q, \quad \mathbf{x} \in \mathcal{Q}, \beta > 0 \quad F(\mathbf{x}) = F_T(\mathbf{x}).$$

- Generate  $\mathbf{x}_t$ 's with  $F_t$ 's and show oracle access of  **$F_t$  is identical to  $F$  in the neighborhood of  $\mathbf{x}_t$**

$$\text{Algorithm } \mathcal{A} : \mathbf{x}_{t+1} = \mathcal{A}(\mathcal{I}_t(\mathbf{x}_1), \dots, \mathcal{I}_t(\mathbf{x}_t)) \quad \text{for} \quad \mathcal{I}_t(\mathbf{x}) = \{F_t, \nabla F_t, \dots, \nabla^p F_t\}.$$

$$F_t(\mathbf{x}) = F(\mathbf{x}) \quad \text{for} \quad \mathbf{x} \in \mathcal{N}_\delta(\mathbf{x}_t), \forall t \in [T]$$

# Hard Function Construction for Case 1: $q > p + \nu$

- Start with a sequence of non-smooth functions

$$g_t(\mathbf{x}) = \max_{1 \leq k \leq t} r_k(\mathbf{x}) \quad \text{where} \quad r_k(\mathbf{x}) = \xi_k \langle \mathbf{e}_{\alpha(k)}, \mathbf{x} \rangle - (k-1)\delta, \forall k \in [T]$$

$\xi_k \in \{-1, 1\}$ ,  $\mathbf{e}$  : standard basis,  $\alpha$  : permutation of  $[T]$ ,  $\delta > 0$ .

- **Smooth the functions** with a smoothing operator, making it  $p^{\text{th}}$ -order **H-smooth**

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

- Add regularization, making it  $q^{\text{th}}$ -order uniformly convex

$$F_t(\mathbf{x}) = \beta G_t(\mathbf{x}) + \frac{\sigma}{q} \|\mathbf{x}\|^q, \quad \mathbf{x} \in \mathcal{Q}, \beta > 0 \quad F(\mathbf{x}) = F_T(\mathbf{x}).$$

- Generate  $\mathbf{x}_t$ 's with  $F_t$ 's and show oracle access of  **$F_t$  is identical to  $F$  in the neighborhood of  $\mathbf{x}_t$**

$$\text{Algorithm } \mathcal{A} : \mathbf{x}_{t+1} = \mathcal{A}(\mathcal{I}_t(\mathbf{x}_1), \dots, \mathcal{I}_t(\mathbf{x}_t)) \quad \text{for} \quad \mathcal{I}_t(\mathbf{x}) = \{F_t, \nabla F_t, \dots, \nabla^p F_t\}.$$

$$F_t(\mathbf{x}) = F(\mathbf{x}) \quad \text{for} \quad \mathbf{x} \in \mathcal{N}_\delta(\mathbf{x}_t), \forall t \in [T]$$

- Lower bound optimality gap  $F(\mathbf{x}_T) - F(\mathbf{x}^*) \geq -\beta(T-1)\delta - \frac{5}{4}p\beta\rho\sqrt{d} + \frac{q-1}{q} \left( \frac{\beta^q}{\sigma T^{\frac{q}{2}}} \right)^{\frac{1}{q-1}}$

## | Smoothing

- Smooth the functions with **a smoothing operator**, making it  $p^{\text{th}}$ -order  $H = \mathcal{O}(1)$ -smooth

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

# | Smoothing

- Smooth the functions with a smoothing operator, making it  $p^{\text{th}}$ -order  $H = \mathcal{O}(1)$ -smooth

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

- Uniform smoothing over a unit  $\ell_2$  ball:  $H = \mathcal{O}(\sqrt{d}) \Rightarrow$  suboptimal rate [Agarwal & Hazan, 2018]

$$S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)], \quad \mathbb{P}[V = \mathbf{v}] = \frac{\Gamma\left(\frac{d}{2} + 1\right)}{\pi^{\frac{d}{2}}} \mathbb{I}_{[\|\mathbf{v}\|_2 \leq 1]}$$

## | Smoothing

- Smooth the functions with a smoothing operator, making it  $p^{\text{th}}$ -order  $H = \mathcal{O}(1)$ -smooth

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

- Uniform smoothing over a unit  $\ell_2$  ball:  $H = \mathcal{O}(\sqrt{d}) \Rightarrow$  suboptimal rate [Agarwal & Hazan, 2018]
- Moreau smoothing:  $H = \mathcal{O}(1)$  [Doikov, 2022], but only 1<sup>st</sup>-order smooth

$$S_\rho[g_t](\mathbf{x}) = \min_{\mathbf{y}} \left\{ g_t(\mathbf{y}) + \frac{\rho}{2} \|\mathbf{y} - \mathbf{x}\|^2 \right\}$$

# | Smoothing

- **Smooth the functions** with a smoothing operator, making it  $p^{\text{th}}$ -order  $H = \mathcal{O}(1)$ -smooth

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

- Uniform smoothing over a unit  $\ell_2$  ball:  $H = \mathcal{O}(\sqrt{d}) \Rightarrow$  suboptimal rate [Agarwal & Hazan, 2018]
- Moreau smoothing:  $H = \mathcal{O}(1)$  [Doikov, 2022], but **only 1<sup>st</sup>-order** smooth
- Softmax smoothing:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth [Bullins, 2020]

$$\text{softmax}_\rho[g_t](\mathbf{x}) = \rho \log \left( \sum_{i=1}^d e^{\frac{x_i}{\rho}} \right)$$

- Gaussian smoothing:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth after  $p$  times [Duchi et. al., 2012]

$$S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)], \quad \mathbb{P}[V = \mathbf{v}] = \frac{1}{(2\pi)^{\frac{d}{2}}} \exp \left\{ -\frac{\mathbf{v}^\top \mathbf{v}}{2} \right\}$$

# | Smoothing

- **Smooth the functions** with a smoothing operator, making it  $p^{\text{th}}$ -order  $H = \mathcal{O}(1)$ -smooth

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

- Uniform smoothing over a unit  $\ell_2$  ball:  $H = \mathcal{O}(\sqrt{d}) \Rightarrow$  suboptimal rate
  - Moreau smoothing:  $H = \mathcal{O}(1)$ , but only 1<sup>st</sup>-order smooth
  - Softmax smoothing:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth
  - Gaussian smoothing:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth after  $p$  times
- Generate  $x_t$ 's with  $F_t$ 's and show oracle access of  $F_t$  is identical to  $F$  in the neighborhood of  $x_t$
- $$F_t(\mathbf{x}) = F(\mathbf{x}) \quad \text{for} \quad \mathbf{x} \in \mathcal{N}_\delta(\mathbf{x}_t), \forall t \in [T]$$
- The smoothing operator needs to be **local**, accessing information only **within the neighborhood**

# Smoothing

- Smooth the functions with a smoothing operator, making it  $p^{\text{th}}$ -order  $H = \mathcal{O}(1)$ -smooth

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

- Uniform smoothing over a unit  $\ell_2$  ball:  $H = \mathcal{O}(\sqrt{d}) \Rightarrow$  suboptimal rate
- Moreau smoothing:  $H = \mathcal{O}(1)$ , but only 1<sup>st</sup>-order smooth
- Softmax smoothing:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth
- Gaussian smoothing:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth after  $p$  times

- Generate  $x_t$ 's with  $F_t$ 's and show oracle access of  $F_t$  is identical to  $F$  in the neighborhood of  $x_t$

$$F_t(\mathbf{x}) = F(\mathbf{x}) \quad \text{for} \quad \mathbf{x} \in \mathcal{N}_\delta(\mathbf{x}_t), \forall t \in [T]$$

- The smoothing operator needs to be **local**, accessing information only **within the neighborhood**
- Gaussian smoothing and softmax smoothing are **global on  $\mathbb{R}^d$**

$$S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$$\mathbb{P}[V = \mathbf{v}] = \frac{1}{(2\pi)^{\frac{d}{2}}} \exp \left\{ -\frac{\mathbf{v}^\top \mathbf{v}}{2} \right\}$$

$$\text{softmax}_\rho[g_t](\mathbf{x}) = \rho \log \left( \sum_{i=1}^d e^{\frac{x_i}{\rho}} \right)$$

# | Smoothing

- **Smooth the functions** with a smoothing operator, making it  $p^{\text{th}}$ -order  $H = \mathcal{O}(1)$ -smooth

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

- Uniform smoothing over a unit  $\ell_2$  ball:  $H = \mathcal{O}(\sqrt{d}) \Rightarrow$  suboptimal rate
  - Moreau smoothing:  $H = \mathcal{O}(1)$ , but only 1<sup>st</sup>-order smooth
  - Softmax smoothing:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth
  - Gaussian smoothing:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth after  $p$  times
- Generate  $x_t$ 's with  $F_t$ 's and show oracle access of  $F_t$  is identical to  $F$  in the neighborhood of  $x_t$
- $$F_t(\mathbf{x}) = F(\mathbf{x}) \quad \text{for} \quad \mathbf{x} \in \mathcal{N}_\delta(\mathbf{x}_t), \forall t \in [T]$$
- The smoothing operator needs to be local, accessing information only within the neighborhood
  - Gaussian smoothing and softmax smoothing are global on  $\mathbb{R}^d$

- **Truncated Gaussian Smoothing**

# | Smoothing

- **Smooth the functions** with a smoothing operator, making it  $p^{\text{th}}$ -order  $H = \mathcal{O}(1)$ -smooth

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

- Uniform smoothing over a unit  $\ell_2$  ball:  $H = \mathcal{O}(\sqrt{d}) \Rightarrow$  suboptimal rate
- Moreau smoothing:  $H = \mathcal{O}(1)$ , but only 1<sup>st</sup>-order smooth
- Softmax smoothing:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth
- Gaussian smoothing:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth after  $p$  times

- Generate  $x_t$ 's with  $F_t$ 's and show oracle access of  $F_t$  is identical to  $F$  in the neighborhood of  $x_t$

$$F_t(\mathbf{x}) = F(\mathbf{x}) \quad \text{for} \quad \mathbf{x} \in \mathcal{N}_\delta(\mathbf{x}_t), \forall t \in [T]$$

- The smoothing operator needs to be local, accessing information only within the neighborhood
- Gaussian smoothing and softmax smoothing are global on  $\mathbb{R}^d$

- **Truncated Gaussian Smoothing**

- Truncated in a unit  $\ell_2$  ball:  $H = \mathcal{O}(\text{poly}(d)) \Rightarrow$  suboptimal rate

$$\mathbb{P}[V = \mathbf{v}] = \frac{1}{Z(d)(2\pi)^{\frac{d}{2}}} \exp \left\{ -\frac{\mathbf{v}^\top \mathbf{v}}{2} \right\} \mathbb{I}_{[\|\mathbf{v}\|_2 \leq 1]}$$

# Smoothing

- **Smooth the functions** with a smoothing operator, making it  $p^{\text{th}}$ -order  $H = \mathcal{O}(1)$ -smooth

$$G_t(\mathbf{x}) = S_\rho^p[g_t](\mathbf{x}) \quad \text{where} \quad S_\rho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} + \rho V)],$$

$V$ : random variable,  $S_\rho^p[\cdot]$ :  $S_\rho[\cdot]$  for  $p$  times,  $\rho > 0$ .

- Uniform smoothing over a unit  $\ell_2$  ball:  $H = \mathcal{O}(\sqrt{d}) \Rightarrow$  suboptimal rate
- Moreau smoothing:  $H = \mathcal{O}(1)$ , but only 1<sup>st</sup>-order smooth
- Softmax smoothing:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth
- Gaussian smoothing:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth after  $p$  times

- Generate  $x_t$ 's with  $F_t$ 's and show oracle access of  $F_t$  is identical to  $F$  in the neighborhood of  $x_t$

$$F_t(\mathbf{x}) = F(\mathbf{x}) \quad \text{for} \quad \mathbf{x} \in \mathcal{N}_\delta(\mathbf{x}_t), \forall t \in [T]$$

- The smoothing operator needs to be local, accessing information only within the neighborhood
- Gaussian smoothing and softmax smoothing are global on  $\mathbb{R}^d$

- **Truncated Gaussian Smoothing**

- Truncated in a unit  $\ell_2$  ball:  $H = \mathcal{O}(\text{poly}(d)) \Rightarrow$  suboptimal rate
- ✓ Truncated in a unit  $\ell_\infty$  ball:  $H = \mathcal{O}(1)$ ,  $p^{\text{th}}$ -order smooth after  $p$  times [Ours, Def. 4 & Lemma 1]

$$\mathbb{P}[V = \mathbf{v}] = \frac{1}{[\Phi(1) - \Phi(-1)]^d (2\pi)^{\frac{d}{2}}} \exp \left\{ -\frac{\mathbf{v}^\top \mathbf{v}}{2} \right\} \mathbb{I}_{\|\mathbf{v}\|_\infty \leq 1}$$

# Hard Function Construction for Case 2: $q < p + \nu$

Orthogonal basis:  $\mathbf{v}_i \perp \mathbf{x}_1, \dots, \mathbf{x}_i$  and  $\mathbf{v}_1, \dots, \mathbf{v}_{i-1}, \forall i \in [\tilde{T}]$

[Arjevani et. al., 2019]  $q = 2, p = 2, \nu = 1$

$$f(\mathbf{x}) = \frac{H}{12} \left( \frac{1}{3} \sum_{i=1}^{\tilde{T}} |\langle \mathbf{v}_i, \mathbf{x} \rangle - \langle \mathbf{v}_{i+1}, \mathbf{x} \rangle|^3 - \gamma \langle \mathbf{v}_1, \mathbf{x} \rangle \right) + \frac{\sigma}{2} \|\mathbf{x}\|^2$$

[Our Work]

$$f(\mathbf{x}) = \frac{H}{2^{p+\nu+1}(p+\nu-1)!} \left( \frac{1}{p+\nu} \sum_{i=1}^{\tilde{T}} |\langle \mathbf{v}_i, \mathbf{x} \rangle - \langle \mathbf{v}_{i+1}, \mathbf{x} \rangle|^{p+\nu} - \gamma \langle \mathbf{v}_1, \mathbf{x} \rangle \right) + \frac{\sigma}{q} \|\mathbf{x}\|^q$$

## Upper Bounds [Song et. al. 2021]

$$q > p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-v)}{q(3(p+v)-2)}} \right)$$

$p = 1$  [Roulet & d'Aspremont, 2017]

$$q = p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \log \left( \frac{1}{\epsilon} \right) \right)$$

[Gasnikov et. al., 2019]

$$q < p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} + \log \log \left( \left( \frac{\sigma^{p+v}}{H^q} \right)^{\frac{1}{p+v-q}} \frac{1}{\epsilon} \right) \right)$$

$q = 2, p = 2, v = 1$  [Arjevani et al., 2019]

## Tight Lower Bounds [Our Work]

$$\Omega \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-v)}{q(3(p+v)-2)}} \right)$$

$p = 1$  [Thomsen & Doikov, 2024]

Future Work

$$\Omega \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} + \log \log \left( \left( \frac{\sigma^{p+v}}{H^q} \right)^{\frac{1}{p+v-q}} \frac{1}{\epsilon} \right) \right)$$

$q = 2, p = 2, v = 1$  [Arjevani et al., 2019]  
 $v = 1$  [Kornowski & Shamir 2020]

## Upper Bounds [Song et. al. 2021]

$$q > p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-v)}{q(3(p+v)-2)}} \right)$$

$p = 1$  [Roulet & d'Aspremont, 2017]

$$q = p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \log \left( \frac{1}{\epsilon} \right) \right)$$

[Gasnikov et. al., 2019]

$$q < p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} + \log \log \left( \left( \frac{\sigma^{p+v}}{H^q} \right)^{\frac{1}{p+v-q}} \frac{1}{\epsilon} \right) \right)$$

$q = 2, p = 2, v = 1$  [Arjevani et al., 2019]

## Tight Lower Bounds [Our Work]

$$\Omega \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-v)}{q(3(p+v)-2)}} \right)$$

$p = 1$  [Thomsen & Doikov, 2024]

Future Work

$$\Omega \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} + \log \log \left( \left( \frac{\sigma^{p+v}}{H^q} \right)^{\frac{1}{p+v-q}} \frac{1}{\epsilon} \right) \right)$$

$q = 2, p = 2, v = 1$  [Arjevani et al., 2019]  
 $v = 1$  [Kornowski & Shamir 2020]

### Ad. #1

On job market, open to research in industry  
& academia, both ML theory & applications.

## Upper Bounds [Song et. al. 2021]

$$q > p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-v)}{q(3(p+v)-2)}} \right)$$

$p = 1$  [Roulet & d'Aspremont, 2017]

$$q = p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \log \left( \frac{1}{\epsilon} \right) \right)$$

[Gasnikov et. al., 2019]

$$q < p + v \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} + \log \log \left( \left( \frac{\sigma^{p+v}}{H^q} \right)^{\frac{1}{p+v-q}} \frac{1}{\epsilon} \right) \right)$$

$q = 2, p = 2, v = 1$  [Arjevani et al., 2019]

## Tight Lower Bounds [Our Work]

$$\Omega \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-v)}{q(3(p+v)-2)}} \right)$$

$p = 1$  [Thomsen & Doikov, 2024]

Future Work

$$\Omega \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+v)-2}} + \log \log \left( \left( \frac{\sigma^{p+v}}{H^q} \right)^{\frac{1}{p+v-q}} \frac{1}{\epsilon} \right) \right)$$

$q = 2, p = 2, v = 1$  [Arjevani et al., 2019]  
 $v = 1$  [Kornowski & Shamir 2020]

### Ad. #1

On job market, open to research in industry  
& academia, both ML theory & applications.

### Ad. #2

Brian is super nice, passionate, knowledgeable.  
You are welcome to collaborate and join us!

## Upper Bounds [Song et. al. 2021]

$$q > p + \nu \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+\nu)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-\nu)}{q(3(p+\nu)-2)}} \right)$$

$p = 1$  [Roulet & d'Aspremont, 2017]

$$q = p + \nu \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+\nu)-2}} \log \left( \frac{1}{\epsilon} \right) \right)$$

[Gasnikov et. al., 2019]

$$q < p + \nu \quad \mathcal{O} \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+\nu)-2}} + \log \log \left( \left( \frac{\sigma^{p+\nu}}{H^q} \right)^{\frac{1}{p+\nu-q}} \frac{1}{\epsilon} \right) \right)$$

$q = 2, p = 2, \nu = 1$  [Arjevani et al., 2019]

## Tight Lower Bounds [Our Work]

$$\Omega \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+\nu)-2}} \left( \frac{\sigma}{\epsilon} \right)^{\frac{2(q-p-\nu)}{q(3(p+\nu)-2)}} \right)$$

$p = 1$  [Thomsen & Doikov, 2024]

Future Work

$$\Omega \left( \left( \frac{H}{\sigma} \right)^{\frac{2}{3(p+\nu)-2}} + \log \log \left( \left( \frac{\sigma^{p+\nu}}{H^q} \right)^{\frac{1}{p+\nu-q}} \frac{1}{\epsilon} \right) \right)$$

$q = 2, p = 2, \nu = 1$  [Arjevani et al., 2019]  
 $\nu = 1$  [Kornowski & Shamir 2020]

**Thank you! Q&A**  
**More @Poster#430 3PM Today**