Tight Lower Bounds under Asymmetric High-Order Hölder Smoothness and Uniform Convexity



Cedar Site Bai bai123@purdue.edu



Brian Bullins
bbullins@purdue.edu



I Oracle Complexity / Convergence Rate

Oracle Complexity

- Objective: $\min_{x} f(x)$
- # times an algorithm accesses an oracle to reach an ϵ -approximate solution

e.g.,
$$f(x_T) - f(x^*) \leq \epsilon$$



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- First-order oracle complexity: # times an algorithm computes gradient



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- First-order Oracle: gradient abla f(x)
- First-order oracle complexity: # times an algorithm computes gradient
- Oracle complexity of gradient descent for convex and smooth functions: $\,T\in\Theta(\epsilon^{-1})\,$

Convergence Rate

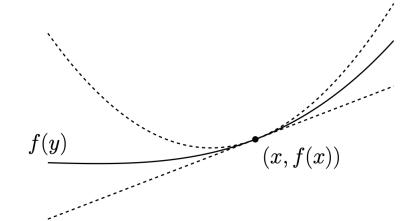
- How fast the approximation error decays with # iterations
- Convergence rate of gradient descent (convex & smooth): $\ f(x_T) f(x^*) \leq \epsilon \in \Theta\left(rac{1}{T}
 ight)$



| Smoothness

$$orall \, \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, L > 0, \|
abla f(\mathbf{x}) -
abla f(\mathbf{y}) \| \leq L \, \| \mathbf{x} - \mathbf{y} \|$$

$$orall \ \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, L > 0, f(\mathbf{y}) - f(\mathbf{x}) - \langle
abla f(\mathbf{x}), \mathbf{y} - \mathbf{x}
angle \leq rac{L}{2} \|\mathbf{x} - \mathbf{y}\|^2$$



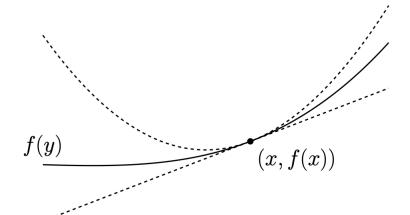


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I Hölder Smoothness

$$egin{aligned} \forall \ \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, H > 0, \pmb{
u} \in (0,1], \|
abla f(\mathbf{x}) -
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u}} \end{aligned}$$





| Smoothness

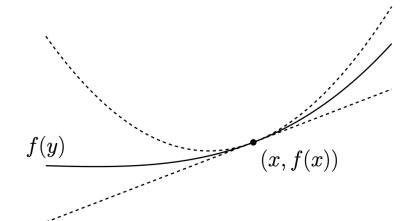
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| High-order Smoothness

$$oldsymbol{p} \in \mathbb{Z}^+, orall \ \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, L > 0, \| oldsymbol{
abla}^p f(\mathbf{x}) - oldsymbol{
abla}^p f(\mathbf{y}) \| \leq L \, \| \mathbf{x} - \mathbf{y} \|$$





| Smoothness

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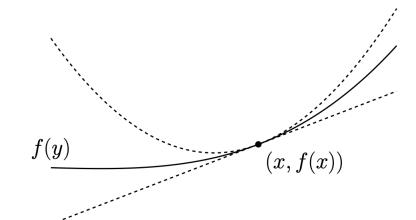
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I Convexity

$$orall \, \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, \langle
abla f(\mathbf{y}) -
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angle \geq 0$$

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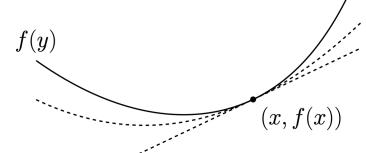


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I Strong Convexity

$$egin{aligned} \forall \ \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, \sigma > 0, \left\langle
abla f(\mathbf{y}) -
abla f(\mathbf{x}), \mathbf{y} - \mathbf{x}
ight
angle \geq \sigma \|\mathbf{x} - \mathbf{y}\|^2 \end{aligned}$$





I Convexity

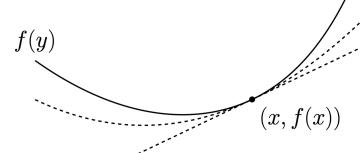
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| Uniform Convexity

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abla f(\mathbf{y}) -
abla f(\mathbf{x}), \mathbf{y} - \mathbf{x}
angle \geq \sigma \|\mathbf{x} - \mathbf{y}\|^{oldsymbol{q}} \end{aligned}$$





Subproblem of Cubic-regularized Newton Method

Let M be a positive parameter. Define a modified Newton step using the following cubic regularization of quadratic approximation of function f(x):

$$T_M(x) \in \text{Arg}\min_{y} \left[\langle f'(x), y - x \rangle + \frac{1}{2} \langle f''(x)(y - x), y - x \rangle + \frac{M}{6} \|y - x\|^3 \right], (2.4)$$

Cubic regularization of Newton method

Initialization: Choose $x_0 \in \mathbb{R}^d$.

Iteration k, $(k \ge 0)$:

- 1. Find $M_k \in [L_0, 2L]$ such that $f(T_{M_k}(x_k)) \leq \bar{f}_{M_k}(x_k)$.
- 2. Set $x_{k+1} = T_{M_k}(x_k)$.



Subproblem of Cubic-regularized Newton Method

$$egin{aligned} F(y) &:= \langle
abla f(x_k), y - x_k
angle + rac{1}{2} \langle
abla^2 f(x_k) (y - x_k), y - x_k
angle + rac{M}{6} \|y - x_k\|^3 \ x_{k+1} &\in rg \min_y F(y) \end{aligned}$$

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- Uniformly convex with degree q=3 and parameter σ :

$$\left| F(y) - F(x) - \langle
abla F(x), y - x
angle \geq rac{\sigma}{q} \|y - x\|^q
ight.$$

- Second-order smooth / Lipschitz Hessian:

$$\left\|
abla^2 F(y) -
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ight\| \leq M \left\| y - x
ight\|$$



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ight\|^{q}$$

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ight\|$$

Ideally, solve with optimal second-order methods [Gasnikov et. al., 2019]

$$p=2, q=3,
u=1, q=p+
u, \implies \mathcal{O}\left(\left(rac{M}{\sigma}
ight)^{rac{2}{7}} \log\left(rac{1}{\epsilon}
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ight)$$



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- Computing Hessian is expensive: $\mathcal{O}(d^2)$
- Use first-order methods, e.g., (accelerated) gradient descent, instead: $\mathcal{O}(d)$

Motivation



| Motivating Example

Subproblem of Cubic-regularized Newton Method

$$F(y) := \langle
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ight)
ight)$$

- Computing Hessian is expensive: $\mathcal{O}(d^2)$
- Use first-order methods, e.g., (accelerated) gradient descent, instead: $\mathcal{O}(d)$
- Order of smoothness limited by order of oracle, e.g., H-smooth (within some domain)

$$p=1, q=3,
u=1, \boxed{q>p+
u} \implies \mathcal{O}\left(\left(\frac{H}{\sigma}\right)^{\frac{1}{2}}\left(\frac{\sigma}{\epsilon}\right)^{\frac{1}{6}}\right)$$
 [Roulet & d'Aspremont, 2017; Song et. al. 2021]



| Uniform Convexity and High-Order Hölder Smoothness

Uniformly Convexity

$$egin{aligned} \forall \ \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, \sigma > 0, q \geq 2, \langle
abla f(\mathbf{y}) -
abla f(\mathbf{x}), \mathbf{y} - \mathbf{x}
angle \geq \sigma \|\mathbf{x} - \mathbf{y}\|^q \end{aligned}$$

• High-order Holder-smoothness

$$p \in \mathbb{Z}^+, orall \, \mathbf{x}, \mathbf{y} \in \mathbb{R}^d, H > 0,
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abla^p f(\mathbf{x}) -
abla^p f(\mathbf{y})\| \leq H \|\mathbf{x} - \mathbf{y}\|^
u$$

Question:

What's the pth-order oracle complexity for different combinations of p, q, ν ?



| Upper Bounds

Upper Bounds [Song et. al. 2021]

$$q>p+
u$$
 $\mathcal{O}\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\left(rac{\sigma}{\epsilon}
ight)^{rac{2(q-p-
u)}{q(3(p+
u)-2)}}
ight)$

$$q = p + \nu \qquad \mathcal{O}\left(\left(\frac{H}{\sigma}\right)^{\frac{2}{3(p+\nu)-2}}\log\left(\frac{1}{\epsilon}\right)\right)$$

$$q$$



| Upper Bounds

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p = 1 [Roulet & d'Aspremont, 2017]

$$q = p + \nu \qquad \mathcal{O}\left(\left(\frac{H}{\sigma}\right)^{\frac{2}{3(p+\nu)-2}}\log\left(\frac{1}{\epsilon}\right)\right)$$

[Gasnikov et. al., 2019]

$$q
$$q = 2, p = 2, \nu = 1 \text{ [Arjevani et al., 2019]}$$$$



Lower Bounds

Upper Bounds [Song et. al. 2021]

Lower Bounds [Our Work]

 $\Omega\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\left(rac{\sigma}{\epsilon}
ight)^{rac{2(q-pu)}{q(3(p+
u)-2)}}
ight)$

Future Work

$$q > p + \nu$$

$$\mathcal{O}\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\left(rac{\sigma}{\epsilon}
ight)^{rac{2(q-p-
u)}{q(3(p+
u)-2)}}
ight)$$

$$p = 1$$
 [Roulet & d'Aspremont, 2017]

$$p=1$$
 [Roulet & d'Aspremont, 2017]

$$q = p + v$$

$$\mathcal{O}\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\log\left(rac{1}{\epsilon}
ight)
ight)$$

$$q
$$q = 2, p = 2, \nu = 1 \text{ [Arjevani et al., 2019]}$$$$

$$\Omega\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}} + \log\log\left(\left(rac{\sigma^{p+
u}}{H^q}
ight)$$



Lower Bounds

Upper Bounds [Song et. al. 2021]

Lower Bounds [Our Work]

$$q > p + \nu$$

$$\mathcal{O}\left(\left(rac{H}{\sigma}
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ight)$$

p = 1 [Thomsen & Doikov, 2024]

$$q = p + \nu$$

$$\mathcal{O}\left(\left(rac{H}{\sigma}
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[Gasnikov et. al., 2019]

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ight)^{rac{1}{p+
u-q}}rac{1}{\epsilon}
ight)
ight)$$
 $q=2, p=2,
u=1$ [Arjevani et al., 2019]
 $u=1$ [Kornowski & Shamir 2020]



| Hard Function Construction for Case 1: $q > p + \nu$

- Start with a sequence of non-smooth functions

$$g_t(\mathbf{x}) = \max_{1 \leq k \leq t} r_k(\mathbf{x}) \quad where \quad r_k(\mathbf{x}) = \xi_k \langle \mathbf{e}_{lpha(k)}, \mathbf{x}
angle - (k-1)\delta, orall \ k \in [T]$$

$$\xi_k \in \{-1,1\}, \ \mathbf{e} : \text{standard basis}, \ \alpha : \text{permutation of } [T], \ \delta > 0.$$



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> Smooth the functions with a smoothing operator, making it pth-order H-smooth

$$G_t(\mathbf{x}) = S^p_
ho[g_t](\mathbf{x}) \quad where \quad S_
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ho[\cdot] \colon S_
ho[\cdot] ext{ for } p ext{ times}, \
ho > 0. \end{aligned}$$

- Add regularization, making it qth-order uniformly convex

$$F_t(\mathbf{x}) = eta G_t(\mathbf{x}) + rac{\sigma}{q} \, \|\, \mathbf{x} \, \|^{\, q}, \quad \mathbf{x} \in \mathcal{Q}, \ eta > 0 \qquad F(\mathbf{x}) = F_T(\mathbf{x}).$$



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- Start with a sequence of non-smooth functions

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 \triangleright Generate \mathbf{x}_t 's with F_t 's and show oracle access of F_t is identical to F in the neighborhood of \mathbf{x}_t

$$egin{aligned} Algorithm \ \mathcal{A}: \ \mathbf{x}_{t+1} &= \mathcal{A}(\mathcal{I}_t(\mathbf{x}_1), \cdots, \mathcal{I}_t(\mathbf{x}_t)) \quad for \quad \mathcal{I}_t(\mathbf{x}) = \{F_t, \nabla F_t, \cdots, \nabla^p F_t\}. \ F_t(\mathbf{x}) &= F(\mathbf{x}) \quad for \quad \mathbf{x} \in \mathcal{N}_\delta(\mathbf{x_t}), orall \ t \in [T] \end{aligned}$$



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$$F_t(\mathbf{x}) = eta G_t(\mathbf{x}) + rac{\sigma}{q} \, ||\, \mathbf{x} \, ||^{\, q}, \quad \mathbf{x} \in \mathcal{Q}, \; eta > 0 \qquad F(\mathbf{x}) = F_T(\mathbf{x}).$$

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$$F_t(\mathbf{x}) = F(\mathbf{x}) \quad for \quad \mathbf{x} \in \mathcal{N}_{\delta}(\mathbf{x_t}), \forall \ t \in [T]$$

 $- \text{ Lower bound optimality gap } F(\mathbf{x}_T) - F(\mathbf{x}^*) \geq -\beta (T-1)\delta - \frac{5}{4}p\beta\rho\sqrt{d} + \frac{q-1}{a}\left(\frac{\beta^q}{2^{q-1}}\right)^{\frac{1}{q-1}}$



| Smoothing

 \triangleright Smooth the functions with **a smoothing operator**, making it pth-order H = $\mathcal{O}(1)$ -smooth

$$G_t(\mathbf{x}) = S^p_
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Department of Computer Science

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V: random variable, $S_{\rho}^{p}[\cdot]$: $S_{\rho}[\cdot]$ for p times, $\rho > 0$.

Uniform smoothing over a unit ℓ_2 ball: $H = \mathcal{O}(\sqrt{d}) \Rightarrow$ suboptimal rate [Agarwal & Hazan, 2018]

$$S_
ho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} +
ho V)], \;\; \mathbb{P}[V = \mathbf{v}] = rac{\Gamma\left(rac{d}{2} + 1
ight)}{\pi^{rac{d}{2}}} \mathbb{I}_{\left[\parallel \mathbf{v} \parallel_2 \leq 1
ight]}$$



I Smoothing

 \triangleright Smooth the functions with a smoothing operator, making it pth-order H = $\mathcal{O}(1)$ -smooth

$$G_t(\mathbf{x}) = S^p_
ho[g_t](\mathbf{x}) \quad where \quad S_
ho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} +
ho V)],$$

- Uniform smoothing over a unit ℓ_2 ball: $H = \mathcal{O}(\sqrt{d}) \Rightarrow$ suboptimal rate [Agarwal & Hazan, 2018]
- Moreau smoothing: $H = \mathcal{O}(1)$ [Doikov, 2022], but only 1st-order smooth

$$S_{
ho}\left[g_{t}
ight]\left(\mathbf{x}
ight)=\min_{\mathbf{y}}\left\{g_{t}(\mathbf{y})+rac{
ho}{2}\|\mathbf{y}-\mathbf{x}\|^{2}
ight\}$$



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ho V)],$$

V: random variable, $S_{\rho}^{p}[\cdot]$: $S_{\rho}[\cdot]$ for p times, $\rho > 0$.

- Uniform smoothing over a unit ℓ_2 ball: $H = \mathcal{O}(\sqrt{d}) \Longrightarrow$ suboptimal rate [Agarwal & Hazan, 2018]
- Moreau smoothing: $H = \mathcal{O}(1)$ [Doikov, 2022], but only 1st-order smooth
- Softmax smoothing: $H = \mathcal{O}(1)$, pth-order smooth [Bullins, 2020]

$$\operatorname{smax}_{
ho}\left[g_{t}
ight]\!\left(\mathbf{x}
ight) =
ho\log\left(\sum_{i=1}^{d}e^{rac{x_{i}}{
ho}}
ight)$$

• Gaussian smoothing: $H = \mathcal{O}(1)$, pth-order smooth after p times [Duchi et. al., 2012]

$$S_
ho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} +
ho V)], \;\; \mathbb{P}\left[V = \mathbf{v}
ight] = rac{1}{(2\pi)^{rac{d}{2}}} \mathrm{exp}\left\{-rac{\mathbf{v}^ op \mathbf{v}}{2}
ight\}$$



I Smoothing

> Smooth the functions with a smoothing operator, making it pth-order H = O(1)-smooth

$$G_t(\mathbf{x}) = S^p_
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- Uniform smoothing over a unit ℓ_2 ball: $H = \mathcal{O}(\sqrt{d}) \Longrightarrow$ suboptimal rate
- Moreau smoothing: $H = \mathcal{O}(1)$, but only 1st-order smooth
- Softmax smoothing: $H = \mathcal{O}(1)$, pth-order smooth
- Gaussian smoothing: $H = \mathcal{O}(1)$, pth-order smooth after p times
- Figure Generate x_t 's with F_t 's and show oracle access of F_t is identical to F in the neighborhood of x_t $F_t(\mathbf{x}) = F(\mathbf{x}) \quad for \quad \mathbf{x} \in \mathcal{N}_{\delta}(\mathbf{x_t}), \forall \ t \in [T]$
 - The smoothing operator needs to be local, accessing information only within the neighborhood



I Smoothing

 \triangleright Smooth the functions with a smoothing operator, making it pth-order H = $\mathcal{O}(1)$ -smooth

$$G_t(\mathbf{x}) = S^p_
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ho[g_t](\mathbf{x}) = \mathbb{E}_V[g_t(\mathbf{x} +
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- Uniform smoothing over a unit ℓ_2 ball: $H = \mathcal{O}(\sqrt{d}) \Longrightarrow$ suboptimal rate
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 - The smoothing operator needs to be local, accessing information only within the neighborhood
 - Gaussian smoothing and softmax smoothing are global on \mathbb{R}^d

$$egin{aligned} S_
ho[g_t](\mathbf{x}) &= \mathbb{E}_V[g_t(\mathbf{x} +
ho V)], \ \mathbb{P}\left[V = \mathbf{v}
ight] &= rac{1}{(2\pi)^{rac{d}{2}}} \mathrm{exp}\left\{-rac{\mathbf{v}^ op \mathbf{v}}{2}
ight\} \end{aligned} \qquad \mathrm{smax}_
ho\left[g_t
ight](\mathbf{x}) =
ho\log\left(\sum_{i=1}^d e^{rac{x_i}{
ho}}
ight).$$



I Smoothing

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$$G_t(\mathbf{x}) = S^p_
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- Truncated Gaussian Smoothing



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 \triangleright Smooth the functions with a smoothing operator, making it pth-order H = $\mathcal{O}(1)$ -smooth

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ho[g_t](\mathbf{x}) \quad where \quad S_
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- Gaussian smoothing and softmax smoothing are global on \mathbb{R}^d
- Truncated Gaussian Smoothing
 - Truncated in a unit ℓ_2 ball: $H = \mathcal{O}(poly(d)) \Longrightarrow$ suboptimal rate

$$\mathbb{P}[V=\mathbf{v}] = rac{1}{Z(d)(2\pi)^{rac{d}{2}}} \mathrm{exp}\left\{-rac{\mathbf{v}^ op \mathbf{v}}{2}
ight\} \mathbb{I}_{\left[\parallel \mathbf{v} \parallel_2 \leq 1
ight]}$$



I Smoothing

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$$F_t(\mathbf{x}) = F(\mathbf{x}) \quad for \quad \mathbf{x} \in \mathcal{N}_{\delta}(\mathbf{x_t}), orall \ t \in [T]$$

- The smoothing operator needs to be local, accessing information only within the neighborhood
- Gaussian smoothing and softmax smoothing are global on \mathbb{R}^d
- > Truncated Gaussian Smoothing
 - Truncated in a unit ℓ_2 ball: $H = \mathcal{O}(poly(d)) \Longrightarrow$ suboptimal rate
 - ✓ Truncated in a unit ℓ_{∞} ball: H = $\mathcal{O}(1)$, pth-order smooth after p times [Ours, Def. 4 & Lemma 1]

$$\mathbb{P}[V=\mathbf{v}] = rac{1}{\left[\Phi(1) - \Phi(-1)
ight]^d (2\pi)^{rac{d}{2}}} \mathrm{exp}\left\{-rac{\mathbf{v}^ op \mathbf{v}}{2}
ight\} \mathbb{I}_{\left[\parallel \mathbf{v} \parallel_{\infty} \leq 1
ight]}$$



| Hard Function Construction for Case 2: q

Orthogonal basis: $\mathbf{v}_i \perp \mathbf{x}_1, \cdots, \mathbf{x}_i \ and \ \mathbf{v}_1, \cdots, \mathbf{v}_{i-1}, \ \forall \ i \in [\tilde{T}]$

[Arjevani et. al., 2019] $q = 2, p = 2, \nu = 1$

$$f(\mathbf{x}) = rac{H}{12} \left(rac{1}{3} \sum_{i=1}^{ ilde{T}} \left| \left\langle \mathbf{v}_i, \, \mathbf{x}
ight
angle - \left\langle \mathbf{v}_{i+1}, \, \mathbf{x}
ight
angle
ight|^3 - \gamma \left\langle \mathbf{v}_1, \, \mathbf{x}
ight
angle
ight) + rac{\sigma}{2} \left\| \, \mathbf{x} \,
ight\|^2$$

[Our Work]

$$f(\mathbf{x}) = rac{H}{2^{p+
u+1}(p+
u-1)!} \left(rac{1}{p+
u} \sum_{i=1}^{ ilde{T}} \left| \left\langle \mathbf{v}_i, \, \mathbf{x}
ight
angle - \left\langle \mathbf{v}_{i+1}, \, \mathbf{x}
ight
angle
ight|^{p+
u} - \gamma \left\langle \mathbf{v}_1, \, \mathbf{x}
ight
angle
ight) + rac{\sigma}{q} \, \| \, \mathbf{x} \, \|^{rac{oldsymbol{q}}{q}}
ight|^{q}$$



Tight Lower Bounds [Our Work]

$$q > p + \nu$$

$$\mathcal{O}\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\left(rac{\sigma}{\epsilon}
ight)^{rac{2(q-p-
u)}{q(3(p+
u)-2)}}
ight)$$

p = 1 [Roulet & d'Aspremont, 2017]

$$\Omega\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\left(rac{\sigma}{\epsilon}
ight)^{rac{2(q-p-
u)}{q(3(p+
u)-2)}}
ight)$$

p = 1 [Thomsen & Doikov, 2024]

Future Work

$$q = p + v$$

$$\mathcal{O}\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\log\left(rac{1}{\epsilon}
ight)
ight)$$

[Gasnikov et. al., 2019]

$$q
$$q = 2, p = 2, \nu = 1 \text{ [Arjevani et al., 2019]}$$

$$q = 2, p = 2, \nu = 1 \text{ [Arjevani et al., 2019]}$$$$

$$q=2, p=2, v=1$$
 [Arjevani et al., 2019]
$$v=1$$
 [Kornowski & Shamir 2020]



Tight Lower Bounds [Our Work]

$$q>p+
u$$
 $\mathcal{O}\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\left(rac{\sigma}{\epsilon}
ight)^{rac{2(q-p-
u)}{q(3(p+
u)-2)}}
ight)$

p = 1 [Roulet & d'Aspremont, 2017]

$$\Omega\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\left(rac{\sigma}{\epsilon}
ight)^{rac{2(q-p-
u)}{q(3(p+
u)-2)}}
ight)$$

p = 1 [Thomsen & Doikov, 2024]

$$q = p + v$$

$$\mathcal{O}\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\log\left(rac{1}{\epsilon}
ight)
ight)$$

[Gasnikov et. al., 2019]

Future Work

$$q
$$q = 2, p = 2, \nu = 1 \text{ [Arjevani et al., 2019]}$$$$

$$\Omega\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}} + \log\log\left(\left(rac{\sigma^{p+
u}}{H^q}
ight)^{rac{1}{p+
u-q}}rac{1}{\epsilon}
ight)
ight)$$
 $q=2,p=2,
u=1$ [Arjevani et al., 2019] $u=1$ [Kornowski & Shamir 2020]

Ad. #1

On job market, open to research in industry & academia, both ML theory & applications.



Tight Lower Bounds [Our Work]

$$q>p+
u$$
 $\mathcal{O}\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\left(rac{\sigma}{\epsilon}
ight)^{rac{2(q-p-
u)}{q(3(p+
u)-2)}}
ight)$

p = 1 [Roulet & d'Aspremont, 2017]

$$\Omega\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\left(rac{\sigma}{\epsilon}
ight)^{rac{2(q-p-
u)}{q(3(p+
u)-2)}}
ight)$$

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ight)^{rac{2}{3(p+
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[Gasnikov et. al., 2019]

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$$q = 2, p = 2, \nu = 1 \text{ [Arjevani et al., 2019]}$$

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$$\Omega\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}} + \log\log\left(\left(rac{\sigma^{p+
u}}{H^q}
ight)^{rac{1}{p+
u-q}}rac{1}{\epsilon}
ight)
ight)$$
 $q=2, p=2,
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 $u=1$ [Kornowski & Shamir 2020]

Ad. #1

On job market, open to research in industry & academia, both ML theory & applications.

Ad. #2

Brian is super nice, passionate, knowledgeable. You are welcome to collaborate and join us!



Tight Lower Bounds [Our Work]

$$q>p+
u$$
 $\mathcal{O}\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\left(rac{\sigma}{\epsilon}
ight)^{rac{2(q-p-
u)}{q(3(p+
u)-2)}}
ight)$

p = 1 [Roulet & d'Aspremont, 2017]

$$\Omega\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\left(rac{\sigma}{\epsilon}
ight)^{rac{2(q-p-
u)}{q(3(p+
u)-2)}}
ight)$$

p=1 [Thomsen & Doikov, 2024]

$$q = p + \nu$$

$$\mathcal{O}\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}}\log\left(rac{1}{\epsilon}
ight)
ight)$$

[Gasnikov et. al., 2019]

Future Work

$$q
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$$\Omega\left(\left(rac{H}{\sigma}
ight)^{rac{2}{3(p+
u)-2}} + \log\log\left(\left(rac{\sigma^{p+
u}}{H^q}
ight)^{rac{1}{p+
u-q}}rac{1}{\epsilon}
ight)
ight)$$
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Thank you! Q&A More @Poster#430 3PM Today